# 14. A Note on Schütte's Interpolation Theorem 

By Nobuyoshi Мотонashi<br>(Comm. by Kunihiko Kodaira, m. J. A., Feb. 12, 1972)

In this note, we shall add some remarks on Schütte's interpolation theorem in the intuitionistic predicate logic (cf. Schütte [3]), one of which give an affirmative solution of one of open problems in Gabbay [1].

Schütte's interpolation theorem. If $A \supset B$ is provable in the intuitionistic predicate logic, then there is a formula $C$ satisfying the following (1) and (2):
(1) $A \supset C$ and $C \supset B$ are provable in this logic.
(2) Every predicate symbol in $C$ occurs both in $A$ and in $B$.

We add the following fact to this theorem:
Theorem. In Schütte's theorem above, if $A$ and $B$ are built up using $\neg$ (negation), $\wedge$ (conjunction) and $\forall$ (universal quantification) only, then we can take such a $C$ which satisfies (1), (2) and an added condition (3) :
(3) Every free variable in $C$ occurs both in $A$ and in $B$.

Remark 1. The proposition obtained from the above theorem by omitting (3) is an affirmative solution of one of open problems in [1].

Remark 2. In Schütte's theorem, we can easily add the condition (3) to C, but in our theorem this is not trivial because we can not apply $\exists$ (existential quantifier) to $C$.

Let $L J$ be the intuitionistic predicate logic formulated by Gentzen in [2]. For the sake of simplicity we assume that a sequent in $L J$ is of the form $\Gamma \rightarrow \Theta$, where $\Gamma$ and $\Theta$ are finite sets of formulas in $L J$ and $\Theta$ has at most one formula, although we shall write $A, \Gamma \rightarrow B$ instead of $\{A\} \cup \Gamma \rightarrow\{B\}$. Furthermore we assume that $L J$ has two propositional constants $T$ (truth), $\perp$ (false) and two added axiom sequents $\rightarrow T$ and $\perp \rightarrow$.

Lemma 1. Let $\Gamma \rightarrow \Theta$ be a sequent in $L J$ and $\left(\Gamma_{1}, \Gamma_{2}\right)$ be an ordered partition of $\Gamma$. If $\vdash_{{ }_{J J}} \Gamma \rightarrow \Theta$, then there is a formula $C$ such that
(4) $\vdash^{{ }_{L J}} \Gamma_{1} \rightarrow C$ and $\vdash^{L J} C, \Gamma_{2} \rightarrow \Theta$.
(5) Every predicate symbol in $C$ occurs both in $\Gamma_{1}$ and $\Gamma_{2} \cup \Theta$.

Furthermore if every formula in $\Gamma \cup \Theta$ is built up using $\neg, \wedge, \forall$ only, then $C$ is also such a formula.
Proof. We use the induction on a cut-free derivation $\mathscr{D}$ of $\Gamma \rightarrow \Theta$. We only treat the case that the last rule of $\mathscr{D}$ is $(\neg \rightarrow)$ or $(\rightarrow \forall)$.

Case 1. The last rule of $\mathscr{D}$ is $(\neg \rightarrow)$. Then $\mathscr{D}$ has the form

