252. An Algebraic Formulation of K-N Propositional Calculus

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A K-N axiom system of propositional calculus is given by J. B. Rosser (2). His axiom system of classical propositional calculus is written in the form of

a) CpKpp,

b) CKpqp,

c) CCpqCNKqrNKrp,

where functors K, N, C denote conjunction, negation, and implication respectively.

As well known, we have Cpq = NKpNq. Therefore Rosser's axiom system is denoted by two functors K, N as follows:

a') NKpNKpp,

b') NKKpqNp,

c') NKNKpNqNNKNKqrNNKrp.

On the other hand, B. Sobociński obtained two new axiom systems which is equivalent to Rosser's system (see B. Sobociński [3], [4]). C. A. Meredith gave an axiom system (see C. A. Meredith and A. N. Prior [1]).

In the K-N propositional calculus, there are two rules of procedure:

1) One of them is the rule of substitution commonly used in the propositional calculus.

2) The other is the rule of detachment as follows. If $NK\alpha N\beta$ and α are theses, then β is also a thesis.

From Rosser's system or KN-system, we can define an algebraic system as follows: Let x be an abstract algebra consisting of $0, p, q, r, \cdots$ with a binary operation * and a unary operation \sim satisfying the following conditions:

1) $\sim (p * p) * p = 0$,

2) $\sim p * (q * p) = 0$,

3) $\sim \sim (\sim \sim (p * r) * \sim (r * q)) * \sim (\sim q * p) = 0$,

4) Let α, β be expressions in X, then $\sim \sim \beta * \sim \alpha = 0$ and $\alpha = 0$ imply $\beta = 0$.

Then X is called KN-algebra. The condition 4) corresponds to the rule of detachment.

First of all, we shall prove some general theorems. The Greek