

A decision method for a set of first order classical formulas and its applications to decision problems for non-classical propositional logics

Dedicated to Professor Shōji Maehara for his sixtieth birthday

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I. Main theorem.

Let L be the first order classical predicate logic without equality. We assume that L has a fixed binary predicate symbol R , unary predicate symbols P_1, \dots, P_N and no other non-logical constant symbols. R -free formulas are formulas in L which has no occurrences of R . R -positive formulas are formulas in L which has no negative occurrences of R . R -formulas are formulas defined inductively as follows:

- (1) All R -free formulas are R -formulas;
- (2) If A and B are R -formulas, then $\neg A, A \wedge B, A \vee B, A \supset B$ are all R -formulas;
- (3) If $A(x)$ is an R -formula and x is a free variable not occurring in $A(v)$, then $\forall v A(v), \forall v(R(x, v) \supset A(v)), \forall v(R(v, x) \supset A(v)), \exists v A(v), \exists v(R(x, v) \wedge A(v)), \exists v(R(v, x) \wedge A(v))$ are all R -formulas.

By R -quantifiers, we denote the quantifiers of the form:

$$\begin{aligned} \forall v(R(x, v) \supset \dots v \dots), & \quad \forall v(R(v, x) \supset \dots v \dots), \\ \exists v(R(x, v) \wedge \dots v \dots), & \quad \exists v(R(v, x) \wedge \dots v \dots), \end{aligned}$$

where $\dots v \dots$ has no occurrences of the free variable x . Then, R -formulas are formulas obtained from R -free formulas by applying propositional connectives, quantifiers and R -quantifiers.

For each R -formula A , let $R\text{-deg}(A)$ be the non-negative integer, called the R -degree of A , defined as follows:

- (1) $R\text{-deg}(A) = 0$ if A is R -free.
- (2) $R\text{-deg}(\neg A) = R\text{-deg}(A)$,
- $R\text{-deg}(A \wedge B) = R\text{-deg}(A \vee B) = R\text{-deg}(A \supset B) = \max\{R\text{-deg}(A), R\text{-deg}(B)\}$,
- (3) $R\text{-deg}(\forall v A(v)) = R\text{-deg}(\exists v A(v)) = R\text{-deg}(A(x))$, and