

A proof of cut-elimination theorem in simple type-theory

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In [4], G. Takeuti conjectured that the cut-elimination theorem would hold in his system GLC as well as in LK. Many attempts to prove it constructively have not yet succeeded. On the other hand, W. Tait [3] proved the cut-elimination theorem for the second order predicate logic by a non-constructive method. In this paper, we shall prove the cut-elimination theorem in simple type-theory also by a non-constructive method. Our proof will be formalizable in Zermelo's set theory, which contains neither the axiom of replacement nor the axiom of choice¹⁾. The author wishes to express his thanks to Professor T. Nishimura, Mr. K. Namba and Mr. T. Uesu for their kind advice and assistance.

§1. Complexes

The system of simple type-theory we shall use is Schütte's system in [2]²⁾. We shall use the notations in [2].

Let V be a semi-valuation³⁾. We shall define V -complexes of type τ by induction on types.

1.1. A V -complex of type 0 is a pair $[e^0, 0]$, where e^0 is an expression of type 0.

1.2. A V -complex of type 1 is a pair $[A, p]$, where A is a well-formed formula and p is t or f satisfying the following conditions.

1.2.1. If A is t in the semi-valuation V , then $p=t$.

1.2.2. If A is f in V , then $p=f$.

1.3. Suppose that the V -complexes of type $\tau_1, \dots, \tau_{n-1}$ and τ_n are already defined. Let $\mathfrak{C}\tau_1, \dots, \mathfrak{C}\tau_n$ be the sets of all the V -complexes of type τ_1, \dots, τ_n

1) Cf. Appendix 2.

2) For the sake of brevity, constants (except function constants) are omitted, since they can be identified with free variables.

3) Our proof remains valid, if the term "semi-valuation" is replaced by "partial valuation" throughout this paper. But we use only the conditions 6.1.1.-6.1.7. in [2] but do not use 6.2.1.-6.2.7. in [2].