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## CORRIGENDUM TO "STRONG NORMALIZATION PROOF WITH CPS-TRANSLATION FOR SECOND ORDER CLASSICAL NATURAL DEDUCTION"

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Our paper [1] contains a serious error. Proposition 4.6 of [1] is actually false and hence our strong normalization proof does not work for the Curry-style  $\lambda\mu$ -calculus. However, our method still can show that (1) the correction of Proposition 5.4 of [2], and (2) the correction of the proof of strong normalization of Church-style  $\lambda\mu$ -calculus by CPS-translation.

Firstly, our method is still effective for the correction of Proposition 5.4 of [2]. The proposition claims that for any Curry-style  $\lambda\mu$ -term u, which is not necessarily typable, if  $u^*$  is strongly normalizable, then u is strongly normalizable too. But its proof does not work, since Proposition 5.1 (i) of [2] is false because of erasing-continuation. Our method proves the similar result for the Curry-style  $\lambda\mu$ -calculus by Propositions 4.3 and 4.12 of [1].

**PROPOSITION.** For any Curry-style  $\lambda \mu$ -term u, if there exists an augmentation  $u^+$  of u such that  $u^{+*}$  is strongly normalizable, then u is strongly normalizable.

Secondly, as mentioned in the concluding remarks of [1], our method is effective for the strong normalization proof of the Church-style  $\lambda\mu$ -calculus, which is called the second-order typed  $\lambda\mu$ -calculus in [2]. The strong normalization of the typed  $\lambda\mu$ -calculus is proved in [2], but its proof with CPS-translation does not work since Proposition 5.5 of [2] is false because of erasing-continuation.

For the Church-style system, the CPS-translation preserves typability of terms, and the strong normalization is proved by our method in [1]. Definition 4.7 in [1] is naturally changed for Church-style terms as follows:

Aug
$$(\mu \alpha^{A}.t) = \{\mu \alpha^{A}.(\lambda z^{\perp}.t^{+})([\alpha^{A}]c^{\forall X.X}\vec{a}); t^{+} \in Aug(t), z^{\perp} \text{ is a fresh} \lambda \text{-variable and } \vec{a} \text{ is a finite sequence of terms and types}\}$$

Then, similarly to the case of the Curry-style, we can prove the following facts, where  $\triangleright_{\lambda}$ ,  $\triangleright_{\mu}$  and  $\triangleright_{\forall}$  are defined as in [2].

- **LEMMAS.** (1) If  $t: \Gamma \vdash A, \Delta$  is provable in the typed  $\lambda \mu$ -calculus, then there is an augmentation  $t^+$  of t such that  $t^+: \Gamma, (\forall X.X)^c \vdash A, \Delta$ .
- (2) If  $t \triangleright_{\lambda}^{1} u$  and  $t^{+}$  is an augmentation of t, then there exists an augmentation  $u^{+}$  of u such that  $t^{+*} \triangleright^{+} u^{+*}$ .

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