

SEQUENCES OF n -DIAGRAMS

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§1. Introduction. We consider only computable languages, and countable structures, with universe a subset of ω , which we think of as a set of constants. We identify sentences with their Gödel numbers. Thus, for a structure \mathcal{A} , the complete (elementary) diagram, $D^c(\mathcal{A})$, and the atomic diagram, $D(\mathcal{A})$, are subsets of ω . We classify formulas as usual. A formula is both Σ_0 and Π_0 if it is open. For $n > 0$, a formula, in prenex normal form, is Σ_n , or Π_n , if it has n blocks of like quantifiers, beginning with \exists , or \forall . For a formula θ , in prenex normal form, we let $neg(\theta)$ denote the dual formula that is logically equivalent to $\neg\theta$ —if θ is Σ_n , then $neg(\theta)$ is Π_n , and vice versa.

DEFINITION 1.1. *For a structure \mathcal{A} , the n -diagram is*

$$D_n(\mathcal{A}) = D^c(\mathcal{A}) \cap \Sigma_n.$$

We are interested in complexity, which we measure by Turing degree. We denote Turing reducibility by \leq_T , and Turing equivalence by \equiv_T . It is clear that for any structure \mathcal{A} , $D_0(\mathcal{A}) \equiv_T D(\mathcal{A})$. We show that for any \mathcal{A} , there exists $\mathcal{B} \cong \mathcal{A}$ such that $D^c(\mathcal{B}) \equiv_T D(\mathcal{B})$. If \mathcal{A} is an algebraically closed field, a real closed field, or any other structure in which we have effective elimination of quantifiers, then this collapse is “intrinsic”; i.e., it happens in all copies. For models of PA , the collapse is not intrinsic. For the standard model of arithmetic, $\mathcal{N} = (\omega, +, \cdot, S, 0)$, we have $D_n(\mathcal{N}) \equiv_T \emptyset^{(n)}$, uniformly in n . In [8], it is shown that for any model \mathcal{A} of PA , there exists $\mathcal{B} \cong \mathcal{A}$ such that $D_{n+1}(\mathcal{B}) \not\leq_T D_n(\mathcal{B})$.

We first consider the following problem.

PROBLEM 1. *Find syntactic conditions on \mathcal{A} guaranteeing that for some n , for all $\mathcal{B} \cong \mathcal{A}$, $D^c(\mathcal{B}) \equiv_T D_n(\mathcal{B})$. In particular, for $n = 0$, find syntactic conditions guaranteeing that for all $\mathcal{B} \cong \mathcal{A}$, $D^c(\mathcal{B}) \equiv_T D(\mathcal{B})$.*

For structures \mathcal{A} that do not exhibit intrinsic collapse of the complete diagram to the atomic diagram, we consider the sequences $(D_n(\mathcal{B}))_{n \in \omega}$ for $\mathcal{B} \cong \mathcal{A}$. We focus on the corresponding sequences of Turing degrees.

DEFINITION 1.2. (i) *For sets X and Y , Y is c.e. in and above X if Y is c.e. relative to X , and $X \leq_T Y$.*

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