

## SEQUENCES OF $n$ -DIAGRAMS

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**§1. Introduction.** We consider only computable languages, and countable structures, with universe a subset of  $\omega$ , which we think of as a set of constants. We identify sentences with their Gödel numbers. Thus, for a structure  $\mathcal{A}$ , the complete (elementary) diagram,  $D^c(\mathcal{A})$ , and the atomic diagram,  $D(\mathcal{A})$ , are subsets of  $\omega$ . We classify formulas as usual. A formula is both  $\Sigma_0$  and  $\Pi_0$  if it is open. For  $n > 0$ , a formula, in prenex normal form, is  $\Sigma_n$ , or  $\Pi_n$ , if it has  $n$  blocks of like quantifiers, beginning with  $\exists$ , or  $\forall$ . For a formula  $\theta$ , in prenex normal form, we let  $neg(\theta)$  denote the dual formula that is logically equivalent to  $\neg\theta$ —if  $\theta$  is  $\Sigma_n$ , then  $neg(\theta)$  is  $\Pi_n$ , and vice versa.

DEFINITION 1.1. *For a structure  $\mathcal{A}$ , the  $n$ -diagram is*

$$D_n(\mathcal{A}) = D^c(\mathcal{A}) \cap \Sigma_n.$$

We are interested in complexity, which we measure by Turing degree. We denote Turing reducibility by  $\leq_T$ , and Turing equivalence by  $\equiv_T$ . It is clear that for any structure  $\mathcal{A}$ ,  $D_0(\mathcal{A}) \equiv_T D(\mathcal{A})$ . We show that for any  $\mathcal{A}$ , there exists  $\mathcal{B} \cong \mathcal{A}$  such that  $D^c(\mathcal{B}) \equiv_T D(\mathcal{B})$ . If  $\mathcal{A}$  is an algebraically closed field, a real closed field, or any other structure in which we have effective elimination of quantifiers, then this collapse is “intrinsic”; i.e., it happens in all copies. For models of  $PA$ , the collapse is not intrinsic. For the standard model of arithmetic,  $\mathcal{N} = (\omega, +, \cdot, S, 0)$ , we have  $D_n(\mathcal{N}) \equiv_T \emptyset^{(n)}$ , uniformly in  $n$ . In [8], it is shown that for any model  $\mathcal{A}$  of  $PA$ , there exists  $\mathcal{B} \cong \mathcal{A}$  such that  $D_{n+1}(\mathcal{B}) \not\leq_T D_n(\mathcal{B})$ .

We first consider the following problem.

PROBLEM 1. *Find syntactic conditions on  $\mathcal{A}$  guaranteeing that for some  $n$ , for all  $\mathcal{B} \cong \mathcal{A}$ ,  $D^c(\mathcal{B}) \equiv_T D_n(\mathcal{B})$ . In particular, for  $n = 0$ , find syntactic conditions guaranteeing that for all  $\mathcal{B} \cong \mathcal{A}$ ,  $D^c(\mathcal{B}) \equiv_T D(\mathcal{B})$ .*

For structures  $\mathcal{A}$  that do not exhibit intrinsic collapse of the complete diagram to the atomic diagram, we consider the sequences  $(D_n(\mathcal{B}))_{n \in \omega}$  for  $\mathcal{B} \cong \mathcal{A}$ . We focus on the corresponding sequences of Turing degrees.

DEFINITION 1.2. (i) *For sets  $X$  and  $Y$ ,  $Y$  is c.e. in and above  $X$  if  $Y$  is c.e. relative to  $X$ , and  $X \leq_T Y$ .*

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