

Π_1^1 RELATIONS AND PATHS THROUGH \mathcal{O}

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§1. Introduction. When bounds on complexity of some aspect of a structure are preserved under isomorphism, we refer to them as *intrinsic*. Here, building on work of Soskov [34], [33], we give syntactical conditions necessary and sufficient for a relation to be intrinsically Π_1^1 on a structure. We consider some examples of computable structures \mathcal{A} and intrinsically Π_1^1 relations R . We also consider a general family of examples of intrinsically Π_1^1 relations arising in computable structures of maximum Scott rank.

For three of the examples, the maximal well-ordered initial segment in a Harrison ordering, the superatomic part of a Harrison Boolean algebra, and the height-possessing part of a Harrison p -group, we show that the Turing degrees of images of the relation in computable copies of the structure are the same as the Turing degrees of Π_1^1 paths through Kleene's \mathcal{O} . With this as motivation, we investigate the possible degrees of these paths. We show that there is a Π_1^1 path in which \mathcal{O} is not computable. In fact, there is one in which no noncomputable hyperarithmetical set is computable.¹ There are paths that are Turing incomparable, or Turing incomparable over a given hyperarithmetical set. There is a pair of paths whose degrees form a minimal pair. However, there is no path of minimal degree.

In Section 2, we summarize earlier results on intrinsically c.e. and intrinsically Σ_α^0 relations. In Section 3, we rework Soskov's results, and we give our result on intrinsically Π_1^1 relations. In Section 4, we describe the examples. In Section 5, we show that for the well-ordered initial segment of the Harrison ordering and related examples, the degrees of images of the relation in computable copies of the structure match those of Π_1^1 paths through \mathcal{O} . In Section 6, we give results on degrees of paths through \mathcal{O} . In the remainder of the present section, we give some background. Most of this material may be found in the book by Ash and Knight [3].

1.1. Kleene's \mathcal{O} . We give a brief description of Kleene's system of notation for computable ordinals. Further details may be found in [29] or [3]. The system consists of a set \mathcal{O} of notations, together with a partial ordering $<_{\mathcal{O}}$. The ordinal 0 gets notation 1. If a is a notation for α , then 2^a is a notation for $\alpha + 1$. Then $a <_{\mathcal{O}} 2^a$, and also, if $b <_{\mathcal{O}} a$, then $b <_{\mathcal{O}} 2^a$. Suppose α is a limit ordinal. If φ_e is

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¹This provides a new solution to Problem 71 on H. Friedman's list [10].