

IS THE EUCLIDEAN ALGORITHM OPTIMAL AMONG ITS PEERS?

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The Euclidean algorithm on the natural numbers  $\mathbb{N} = \{0, 1, \dots\}$  can be specified succinctly by the *recursive program*

$$\varepsilon : \gcd(a, b) = \begin{cases} b, & \text{if } \text{rem}(a, b) = 0, \\ \gcd(b, \text{rem}(a, b)), & \text{otherwise} \end{cases} \quad (a \geq b \geq 1),$$

where  $\text{rem}(a, b)$  is the remainder in the division of  $a$  by  $b$ , the unique natural number  $r$  such that for some natural number  $q$ ,

$$(1) \quad a = bq + r \quad (0 \leq r < b).$$

It is an algorithm *from (relative to)* the remainder function  $\text{rem}$ , meaning that in computing its *time complexity function*  $c_\varepsilon(a, b)$ , we assume that the values  $\text{rem}(x, y)$  are provided on demand by some “oracle” in one “time unit”. It is easy to prove that

$$c_\varepsilon(a, b) \leq 3 \log_2 a \quad (a > b > 1).$$

Much more is known about  $c_\varepsilon(a, b)$ , but this simple-to-prove upper bound suggests the proper formulation of the Euclidean’s (worst case) optimality among its *peers*—algorithms from  $\text{rem}$ :

**CONJECTURE.** *If an algorithm  $\alpha$  computes  $\gcd(x, y)$  from  $\text{rem}$  with time complexity  $c_\alpha(x, y)$ , then there is a rational number  $r > 0$  such that for infinitely many pairs  $a > b > 1$ ,  $c_\alpha(a, b) > r \log_2 a$ .*

Our main aim here is to prove the following relevant result:

**THEOREM A.** *If a recursive program  $\alpha$  decides the coprimeness relation  $x \perp y$  from  $=, <, +, \div, \text{iq}$  and  $\text{rem}$ , then for infinitely many coprime pairs  $a > b > 1$ ,*

$$(2) \quad c_\alpha(a, b) > \frac{1}{10} \log_2 \log_2 a.$$

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