

## A NOTE ON A FORMALIZED ARITHMETIC WITH FUNCTION SYMBOLS ' AND +.

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### Introduction.

Let  $\mathfrak{L}_0$  be the first order language with function symbols ', + and the equality symbol =. By  $\mathfrak{L}$  we denote the first order language obtained from  $\mathfrak{L}_0$  by adding a ternary predicate symbol  $P$ . The theory in  $\mathfrak{L}$  with the following axioms and axiom schemata is signified by  $\mathfrak{N}$ .

- (N- 1)  $\forall x \neg(x' = 0)$ .
- (N- 2)  $\forall x \forall y (x' = y' \supset x = y)$ .
- (N- 3)  $\forall x (x + 0 = x)$ .
- (N- 4)  $\forall y \forall y' (x + y' = (x + y)')$ .
- (N- 5)  $\forall x P(x, 0, 0)$ .
- (N- 6)  $\forall x \forall y \forall z \{P(x, y, z) \supset P(x, y', z + x)\}$ .
- (N- 7)  $\forall x \forall y \forall z \forall w \{(P(x, y, z) \wedge P(x, y, w)) \supset z = w\}$ .
- (N- 8)  $\forall x (x = x)$ .
- (N- 9)  $\forall x \forall y \{x = y \supset (\mathfrak{A}(x) \supset \mathfrak{A}(y))\}$ .
- (N-10)  $\{\mathfrak{A}(0) \wedge \forall x (\mathfrak{A}(x) \supset \mathfrak{A}(x'))\} \supset \forall x \mathfrak{A}(x)$ .
- (N-11)  $s = t$ , where  $s = t$  is valid.

For a term  $t$ ,  $b(t)$  means the number of occurrences of bound variables in  $t$ . For a formula  $\mathfrak{A}$ ,  $b(\mathfrak{A})$  is defined inductively as follows. 1.  $b(r = s) = \max(b(r), b(s))$ . 2.  $b(P(r, s, t)) = \max(b(r), b(s), b(t))$ . 3.  $b(\neg \mathfrak{A}) = b(\mathfrak{A})$ . 4.  $b(\mathfrak{A} \wedge \mathfrak{B}) = b(\mathfrak{A} \vee \mathfrak{B}) = \max(b(\mathfrak{A}), b(\mathfrak{B}))$ . 5.  $b(\forall x \mathfrak{A}) = b(\exists x \mathfrak{A}) = b(\mathfrak{A})$ .

In [3] we proved that:

*For any formula  $\mathfrak{A}(a)$  of  $\mathfrak{L}$ ; if there is a number  $m$  such that, for any natural number  $n$ , there exists a proof  $\mathfrak{P}$  of  $\mathfrak{A}(\bar{n})$  in  $\mathfrak{N}$  with the following properties (1) and (2), then  $\forall x \mathfrak{A}(x)$  is provable in  $\mathfrak{N}$ .*

- (1) *The length of  $\mathfrak{P}$  is less than  $m$ .*
- (2) *For any induction schema  $\mathfrak{B}$  in  $\mathfrak{P}$  which is not a formula of  $\mathfrak{L}_0$ ,  $b(\mathfrak{B}) \leq m$ .*

The purpose of this paper is to prove the following theorem.

**THEOREM.** *There are a formula  $\mathfrak{A}(a)$  and a natural number  $M$  such that: (a)*