

ANOTHER PROOF OF THE STRONG COMPLETENESS OF THE INTUITIONISTIC FUZZY LOGIC

By

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Takeuti and Titani [3] introduced the system, which we shall call TT, for the intuitionistic fuzzy logic, and proved the following theorem:

STRONG COMPLETENESS THEOREM (Takeuti and Titani [3, Theorem 1.3]).
Suppose that the language of TT is countable. If a sequent $\Sigma \Rightarrow \Delta$ is valid then it is provable in TT, where Σ may be infinite.

The purpose of this note is to give another proof of the above theorem.

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§1. Recall, first, that the axioms and inference rules of TT are those of the intuitionistic logic (Gentzen's LJ) together with the following ones:

EXTRA AXIOM SCHEMATA FOR TT.

1. $\Rightarrow(A \rightarrow B) \vee ((A \rightarrow B) \rightarrow B)$;
2. $(A \rightarrow B) \rightarrow B \Rightarrow (B \rightarrow A) \vee B$;
3. $(A \wedge B) \rightarrow C \Rightarrow (A \rightarrow C) \vee (B \rightarrow C)$;
4. $A \rightarrow (B \vee C) \Rightarrow (A \rightarrow B) \vee (A \rightarrow C)$;
5. $\forall x(C \vee A(x)) \Rightarrow C \vee \forall x A(x)$, where x does not occur in C ;
6. $\forall x A(x) \rightarrow C \Rightarrow \exists x(A(x) \rightarrow D) \vee (D \rightarrow C)$, where x does not occur in D .

EXTRA INFERENCE RULE FOR TT.

$$\frac{\Gamma \Rightarrow A \vee (C \rightarrow p) \vee (p \rightarrow B)}{\Gamma \Rightarrow A \vee (C \rightarrow B)},$$

where p is any propositional variable not occurring in the lower sequent.

We call that system TT^- which is obtained from TT by deleting Extra Inference Rule for TT.