

## AN AUGMENTED MODAL LOGIC

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The purpose of this paper is to examine the extent to which an augmented modal logic may be used as the formalized meta-theory  $\mathfrak{M}$  of a formal system  $\mathcal{S}$ . The conditions which  $\mathcal{S}$  must satisfy are indicated by Roman numerals. Assumptions and theorems of  $\mathfrak{M}$  are indicated by Arabic numerals. Comments and examples are enclosed in square brackets.

I Let  $\mathcal{S}$  consist of the following: (i) a set  $\mathbf{V}$  of symbols,  $v^1, v^2, \dots, v^j, \dots$ , (the vocabulary of  $\mathcal{S}$ ); (ii) formation rules; (iii) axioms; (iv) transformation rules.

II Let  $\mathbf{F}$  be the set of formulae  $f^1, f^2, \dots, f^j, \dots$ , of  $\mathcal{S}$ , where each  $f^i$  consists of a finite string of symbols  $v^i$ .

III Let the following sub-sets of  $\mathbf{F}$  be selected as follows:

- (i) A set  $\mathbf{W}$  of wff selected recursively from  $\mathbf{F}$  by the formation rules.
- (ii) A set  $\mathbf{A}$  (the axioms of  $\mathcal{S}$ ) selected recursively from  $\mathbf{W}$  (commonly by giving a finite list).
- (iii) A set  $\mathbf{P}$  of provable formulae; a set  $\mathbf{D}$  of disprovable formulae, a set  $\mathbf{I}$  of irresolvable formulae; selected (but not necessarily recursively) from  $\mathbf{W}$  by the transformation rules. Let  $\mathbf{P}$  contain  $\mathbf{A}$ .

IV Let the operations required to select the various sub-sets of  $\mathbf{F}$  be entirely formal. [By this we mean that they depend only on the physical characteristics—shape, position, etc.,—of the  $v^i$ ].

V Let the sets  $\mathbf{P}$ ,  $\mathbf{D}$  and  $\mathbf{I}$  satisfy the following conditions:

- (i) If, for a given  $j$ ,  $f^j$  occurs in  $\mathbf{P}$ , then for some definite symbol,  $v^k$  say,  $v^k f^j$  occurs in  $\mathbf{D}$ ; and conversely, if  $f^j$  occurs in  $\mathbf{D}$  then  $v^k f^j$  occurs in  $\mathbf{P}$ . [Commonly  $v^k$  will be the symbol ' $\sim$ ' (interpretable as 'not') and we say that  $f^j$  and  $\bar{v}^k f^j$  is each the negation of the other. We leave it open at this stage whether  $\mathbf{P}$  (or  $\mathbf{D}$ ) can contain both  $f^j$  and  $v^k f^j$ , for a given  $j$ ].
- (ii) If neither  $f^j$  nor  $v^k f^j$  occurs in  $\mathbf{P}$  (in which case, by (i) above, neither occurs in  $\mathbf{D}$ ) then both occurs in  $\mathbf{I}$ .