# ON THE CONNECTION OF THE FIRST-ORDER FUNCTIONAL CALCULUS WITH $\aleph_{0}$ PROPOSITIONAL CALCULUS 

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A simply conclusion from papers [2]-[5] is that for each formula $E$ we may construct a $n(E)$-valued propositional calculus such that if $E$ is not a thesis, then $E$ is false in this calculus by a finite interpretation of the quantifiers; by means of a simply extending of the $n(E)$ valued calculus to $\aleph_{0}$ propositional calculus we may prove in one the converse theorem. This method we have used in [5] and have proved that it is possible to approximate the first-order functional calculus by many valued propositional calculi.

An interest approximation of the first-order functional calculus by $\aleph_{0}$ propositional calculus follows from [3] and [4]. We obtain it by means of constructing of a correspondence between atomic formulas and sequences of numbers 0 and 1 such that:

1. If the atomic formula is of $\geq 2$ arguments, then the correspondents sequence is periodic/we shall give the period/.
2. The difference in this correspondence is in general on atomic formulas of one argument whose we must consider an infinite number.
3. For some formulas, e.g. $\Sigma a_{1} \Sigma a_{2} \Pi a_{3} \ldots \Pi a_{k} F$ where $\bar{F}$ is quantifier and individual variable-free, monadic formulas, ..., the $\aleph_{0}$ calculus may be replaced by suitable $n$ - or 2 -valued propositional calculus; one follows from a general theorem.

We shall use the notation of all mentioned papers and in particular:
(1) variables: ( $1^{\circ}$ ) individual: $x_{1}, x_{2}, \ldots$ /or simply $x /,\left(2^{\circ}\right)$ apparent: $a_{1}, a_{2}, \ldots$ /or simply $a /$,
(2) finite numbers of functional variables: $f_{1}^{1}, \ldots, f_{q}^{1}, f_{1}^{2}, \ldots, f_{q}^{2}, \ldots, f_{1}^{t}$, $\ldots, f_{\bar{q}}^{t} / f_{i}^{m}$ of $m$-arguments, $m=1, \ldots, t$ and $i=1, \ldots q /$
(3) logical constants: (negation), + (alternative), $\Pi$ (general quantifier),
(4) atomic expression: $R, R_{1}, R_{2}, \ldots$; expressions: $E, F, G, E_{1}, F_{1}$, $G_{1} \ldots{ }^{1}$

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[^0]:    1. Expressions and formulas we define in the usual way; the expression in which an apparent variable $a$ belong to the scope of two quantifiers $\Pi a$ is not a formula; if $a$ does not occur in $E$, then $\Pi a E$ is not a formula.
