

# FUNCTIONAL COMPLETENESS OF HENKIN'S PROPOSITIONAL FRAGMENTS

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It was shown in [1] that if  $\phi(x_1, \dots, x_2)$  has the defined property of being Tarskian, the addition of schemata  $(\phi)^*$  as in [2] to the positive logic of implication, *A1-2*, yields the complete system of classical implication. Knowledge of [1] is pre-supposed. We define:

Def.  $\mathfrak{T}_1$  For all  $\phi$ ,  $\phi$  is Tarskian<sub>1</sub> iff  $\phi$  is Tarskian and is valued F when all its arguments are valued T.

Def.  $\mathfrak{T}_2$  For all  $\phi$ ,  $\phi$  is Tarskian<sub>2</sub> iff  $\phi$  is Tarskian and is valued T when all its arguments are valued T.

THEOREM 1. If  $\phi$  is Tarskian<sub>1</sub>,  $\{A1-2, (\phi)^*\}$  is functionally complete.

Proof. If  $\phi$  is Tarskian<sub>1</sub>, the proof of Lemma 1, case (i)a and the corresponding sub-case of case (ii), in [1], shows that  $\{A1-2, (\phi)^*\}$  contains *S1-2* with each *i*-th argument of  $\phi$  either *A* or  $A \supset A$ . Defining the negation of *A* for  $\phi$  with these arguments, we get from *S1-2*:

$$(1) A \supset . \sim A \supset C$$

$$(2) A \supset C \supset . \sim A \supset C \supset C.$$

Taking *C* in (2) as *A*, and detaching  $A \supset A$ , we get

$$(3) \sim A \supset A \supset A.$$

Since hypothetical syllogism is given by *A1-2*, and this with (1) and (3) constitutes the well known Łukasiewicz base for a full and functionally complete system in implication and negation, the theorem follows.

THEOREM 2. If  $\phi$  is valued T when all its arguments are valued T, negation is not definable in the system  $\{A1-3, (\phi)^*\}$ .

Proof. The system  $\{A1-3, (\phi)^*\}$  is, by [2], complete for tautologies in implication and  $\phi$ . So every expression  $A \supset B$  with *A* and *B* tautologous is provable, and by the hypothesis on  $\phi$ ,  $\phi(A \supset A, \dots, A \supset A)$  is provable. Hence every expression *f* (imp,  $\phi$ ,  $A \supset A$ ) with implication and  $\phi$  as the only functors, and all elementary argument places filled by  $A \supset A$ , is provable. We suppose now that negation is definable. We should have as provable