

## LEŚNIEWSKI'S ONTOLOGY AND SOME MEDIEVAL LOGICIANS

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In the issue of this journal dated October 1966 (Vol. VII, No. 4, pp. 361-364) Professor John Trentman suggested limitations on my claim that Leśniewski's Ontology is of use in furnishing formal analyses of medieval logical theories, his grounds being that certain medieval theories deny what is called the "two-name theory of predication" allegedly common to Ockham and Ontology. Hence while the work of Ockhamists would be analysable with reference to Ontology, that of those "Thomists" who deny the two-name theory would not. Professor Trentman then goes on to suggest that for such "Thomist" analyses to take place, "something like Frege's functional analysis of predication", with a form like " $\phi(A)$ " is needed to show the "disparity of semantic category that holds between the subject and the predicate", thereby implying that no such form is available in Ontology, and that the allegations about the inadequacy of the two-name theory could have escaped my notice.

Neither of these implications is tenable. Ignoring the second of them, I can deal with the first by exemplifying the manner in which the Ontology in question deals with the relations between names and verbs (i.e. functors which when completed with nominal arguments form propositions). Thus definitions S6.21.12 and S6.22.11 from *The De Grammatico of St. Anselm* (D. P. Henry, Notre Dame 1964) run as follows:

- 1)  $[ab]: \varepsilon\{b\}(a) . \equiv . a \varepsilon b$
- 2)  $[a\phi]: a \varepsilon \text{trm} \langle \phi \rangle . \equiv . a \varepsilon a . \phi(a)$

From these one may infer:

- 3)  $[a]: . a \varepsilon a . \supset: [\exists \phi] . \phi(a) . \equiv . [\exists b] . a \varepsilon b$

The functor defined by 1) is ' $\{ \}$ ', a functor-forming functor for one argument which is a name, the functor thus formed, when completed with one nominal argument, yielding a proposition; it is thus one of many instances of the ' $\phi$ ' of Professor Trentman's preferred form which Ontology makes available, and guarantees a verb corresponding to every name, and hence a ' $\phi(A)$ ' form corresponding to every 'two-name' form of the type of ' $a \varepsilon b$ '.