

GENERALIZATIONS OF THE DISTRIBUTIVE  
 AND ASSOCIATIVE LAWS

ALAN C. WILDE

1 Introduction Let  $x \Delta y$  and  $x \circ y$  denote two truth-value functions:  $\{0, 1\} \times \{0, 1\} \rightarrow \{0, 1\}$ , where 1 and 0 denote "true" and "false" respectively. The two functions "and" and "or" satisfy the law

$$(*) \quad x \Delta (y \circ z) = (x \Delta y) \circ (x \Delta z)$$

in either order. We would like to weaken (\*) so that more functions satisfy the relationship. To do so, we use

$$(**) \quad x \Delta (y \circ z) = (x \Delta y) \circ (x \Delta z) \circ (x \Delta I)$$

where I is the identity of  $x \circ y$ . (\*\*) is a generalization of (\*) for the reason that all functions  $x \circ y$  that have identities and all  $x \Delta y$  that together satisfy (\*) also satisfy (\*\*), but not conversely. This is shown in Theorem 1.

"And" and "or" satisfy the associative law

$$x \Delta (y \Delta z) = (x \Delta y) \Delta z,$$

and so does "equivalence" and "exclusive or." However, we shall demonstrate that for all truth-functions  $x \Delta y$ , the truth-values of  $x \Delta (x \Delta z) \equiv (x \Delta y) \Delta z$  and  $x \Delta (y \Delta z) \underline{\vee} (x \Delta y) \Delta z$  are independent of  $y$ .

2 The Generalized Distributive Law We wish to prove the following:

Theorem 1 (\*\*) holds

(a) for all  $x \Delta y$  if  $x \circ y$  is either  $x \equiv y$  or  $x \underline{\vee} y$ ;

and

(b) for all  $x \Delta y$  such that  $y \leq z$  implies  $x \Delta y \leq x \Delta z$  if  $x \circ y$  is  $x \underline{\vee} y$  or  $x \wedge y$ .

*Proof:* Note that  $x \wedge y$ ,  $x \underline{\vee} y$ ,  $x \equiv y$ , and  $x \underline{\vee} y$  are the only functions that have identities, so Theorem 1 has all the possible combinations. All four of them happen to be commutative and associative. For part (a), let us show