

## A NEW REPRESENTATION OF S5

STEVEN K. THOMASON

We consider first a modal language with propositional constants (and no variables) and show that there is a unique set  $H$  of formulas of this language meeting certain attractive syntactical conditions; moreover  $H$  is the set of theses of a very simple calculus. We then show that the theses of S5 are characterized by the fact that all their instances are in  $H$ .\*

Let  $\mathcal{L}_c$  be the language having an infinite set of "propositional constants" and connectives  $\neg$ ,  $\vee$ , and  $\Box$  used in the usual way. As usual, other connectives are used as abbreviations. If  $S$  is a string of symbols,  $s_1, \dots, s_n$  are distinct symbols, and  $S_1, \dots, S_n$  are strings of symbols, then  $S(S_1, \dots, S_n/s_1, \dots, s_n)$  is the result of replacing each symbol  $s_i$  ( $i = 1, \dots, n$ ) in  $S$  by the string  $S_i$ . A *tautology* is a string of the form  $X(S_1, \dots, S_n/x_1, \dots, x_n)$  where  $X$  is a tautology of the classical propositional calculus and  $x_1, \dots, x_n$  are propositional variables. A set  $H$  of formulas of  $\mathcal{L}_c$  is *correct* if for all formulas  $A$  and  $B$  of  $\mathcal{L}_c$

- (1) If  $A$  is a tautology then  $A \in H$ .
- (2) If  $A$  has no occurrences of  $\Box$  and  $A \in H$ , then  $A$  is a tautology.
- (3) If  $A \in H$  and  $A \Rightarrow B \in H$ , then  $B \in H$ .
- (4)  $A \in H$  if and only if  $\Box A \in H$ .
- (5) Either  $A \in H$  or  $\neg \Box A \in H$ .

Let  $\mathcal{L}_v$  be the language which is like  $\mathcal{L}_c$  except that  $\mathcal{L}_v$  has a countably infinite set of "propositional variables" rather than propositional constants. A set  $J$  of formulas of  $\mathcal{L}_v$  is said to be correct if it consists of all formulas  $X$  of  $\mathcal{L}_v$  such that every formula of  $\mathcal{L}_c$  of the form  $X(A_1, \dots, A_n/x_1, \dots, x_n)$  is a member of  $H$ , where  $H$  is a correct set of formulas of  $\mathcal{L}_c$ .

Let  $\mathfrak{C}$  be the formal system whose language is  $\mathcal{L}_c$ , whose axioms are an appropriate set of tautologies and all formulas of the form

$$\Diamond \& \{a_i^* \mid i = 1, \dots, n\}$$

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