Notre Dame Journal of Formal Logic
Volume XIV, Number 1, January 1973
NDJFAM

EFFECTIVE EXTENDABILITY AND FIXED POINTS

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Let α be any sequence and let $\varphi_1, \varphi_2, \ldots$ be a standard enumeration of the partial recursive functions. A p.r.f. δ is said to be a *fixed-point algorithm* for α if and only if $\delta(n)$ is an α-fixed point for φ_n (i.e., $n \in \mathrm{Dom}\,\delta$ and $\alpha(\delta(n)) = \alpha(\varphi_n(\delta(n)))$ whenever φ_n is total). α has the *effective fixed-point property* if and only if α has a total fixed-point algorithm. The purpose of this paper is to show that the effective fixed-point property is more properly viewed as an extendability property since:

(1) *α has the* e.f.p.p. *if and only if every partial recursive function ψ has a total recursive α-extension f* (i.e., $\alpha(f(n)) = \alpha(\psi(n))$ *for all* $n \in \mathrm{Dom}\,\psi$).
(2) *There is a sequence having a fixed-point algorithm but not the* e.f.p.p. (Hence totalness of the fixed-point algorithm is crucial to the e.f.p.p.)
(3) *If there is a total recursive function f such that $f(x)$ is an α-fixed point of φ_x whenever φ_x is total and constant, then α has the* e.f.p.p. (Hence the fixed points are somewhat incidental to the e.f.p.p. since every sequence has a nontotal algorithm which finds fixed points for constant functions, for example, $\lambda x[\varphi_x(1)]$.)

Proof of 1. See [3], Lemma 1.1.
Proof of 2. We let α be the canonical sequence of equivalence classes associated with the equivalence relation $\approx$ constructed below. Along with $\approx$ we construct a partial recursive function ψ having no total recursive α-extension. Thus α lacks the e.f.p.p. by (1).

Let $T_1, T_2, \ldots$ be a recursive sequence of disjoint infinite recursive sets. Members of T_x are called *test values* for φ_x. Let f be a one to one recursive enumeration of $\{\langle x, y\rangle \mid y \in \mathrm{Dom}\,\varphi_x\}$. We suppose that $\varphi_1, \varphi_2, \ldots$ are being constructed in stages so that $\varphi_x(y)$ becomes defined at stage $f^{-1}(\langle x, y\rangle)$ and at this stage we perform the following three steps in the construction of $\approx$ and ψ:

(Step 1) If φ_x does not already have an α-fixed point we give it one by letting $y \approx \varphi_x(y)$ provided that we do not thereby cause the violation of a prohibition of order x or less.

(Step 2) If y is a test value of φ_x and φ_x agrees, modulo $\approx$, with ψ

Received December 2, 1972