

AXIOMATIC, SEQUENZEN-KALKUL, AND SUBORDINATE
 PROOF VERSIONS OF S9

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1.1 *The System S9.* In [8] the system S9 was presented with primitive connectives \sim , $\&$, and \rightarrow as S3 plus the axioms

- (a) $\sim p \vee ((\sim(p \& \sim p) \rightarrow p) \vee (p \rightarrow (p \rightarrow p)))$,
 (b) $\sim p \vee ((p \rightarrow p) \rightarrow p)$,
 (c) $(p \rightarrow p) \rightarrow \sim(\sim(p \rightarrow p) \rightarrow (p \rightarrow p))$,

the rules being Substitution, Strict Detachment, and Adjunction.¹ A simpler formulation of S9 can be given, however, in that (b) is redundant and (a) and (c) can be simplified. If we abbreviate $\sim x \rightarrow x$ by $\Box x$ and $\sim(x \& \sim y)$ by $x \supset y$, then in S3, $x \rightarrow y$ is strictly equivalent (s.e.) to $\Box(x \supset y)$ and $x \rightarrow x$ is s.e. to $\Box t$, where t is any tautology of classical two-valued logic, PC. Thus, in S3 the axioms (a), (b), and (c) are s.e. to

- (d) $\sim(p \rightarrow \Box t) \supset (p \supset \Box p)$,
 (e) $p \supset (\Box t \rightarrow p)$,
 (f) $\Box t \rightarrow \sim \Box \Box t$

respectively. Now (f) is derivable from $\sim \Box \Box t$ and (e), and (e) is derivable from $\sim \Box \Box t$ and (d). The latter is shown as follows. The formula $(\sim p \rightarrow q) \& (p \rightarrow q) \rightarrow \Box q$ is provable in S3, so that $\sim \Box \Box t \rightarrow ((\sim p \rightarrow \Box t) \supset \sim(p \rightarrow \Box t))$ is provable in S3. Hence, by $\sim \Box \Box t$ and (d) we have $(\sim p \rightarrow \Box t) \supset (p \supset \Box p)$. Substituting $\Box t \supset p$ for p and detaching $\sim(\Box t \supset p) \rightarrow \Box t$, we have $(\Box t \supset p) \supset (\Box t \rightarrow p)$, which yields (e) by a two-valued tautology. Hence, in S3 (a), (b), and (c) are derivable from $\sim \Box \Box t$ and $\sim(p \rightarrow \Box t) \supset (p \supset \Box p)$, and *vice versa*. A simpler formulation of S9 in \sim , $\&$, and \rightarrow is thus S3 plus

- (g) $\sim(\sim(p \rightarrow p) \rightarrow (p \rightarrow p))$
 (h) $\sim(p \rightarrow (p \rightarrow p)) \supset (p \supset (\sim p \rightarrow p))$.

1.2 It is desirable to present yet another formulation of S9: a *Lemmon*

1. For a detailed discussion of S9 see [8].