

SIMPLIFIED FORMALIZATIONS OF FRAGMENTS OF THE PROPOSITIONAL CALCULUS

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Henkin has given [1] a general method of formalizing 2-valued propositional calculi whose primitive functors are such that material implication is definable in terms of them. Let the primitive functors, other than implication if implication is a primitive functor, be the functors F_i of n_i arguments ($i = 1, \dots, b$) and let the formulae $P_1, \dots, P_{n_i}, F_i P_1 \dots P_{n_i}$ take the truth-values $x_1, \dots, x_{n_i}, f_i(x_1, \dots, x_{n_i})$ respectively ($i = 1, \dots, b$). The axiom schemes are

A1 $CPCQP$,

A2 $CCPQCCPCQRCPR$,

A3 $CCPRCCCPQRR$,

A4 $CV_{x_1}P_1Q \dots CV_{x_{n_i}}P_{n_i}QV_yF_iP_1 \dots P_{n_i}Q(y = f_i(x_1, \dots, x_{n_i});$
 $x_1 = \mathbf{T}, \mathbf{F}; \dots; x_{n_i} = \mathbf{T}, \mathbf{F}; i = 1, \dots, b)$,

A4 denoting $\sum_{i=1}^b 2^{n_i}$ axiom schemes and the functors $V_{\mathbf{T}}, V_{\mathbf{F}}$ being defined by the equations

$$V_{\mathbf{T}}PQ =_{df} CCPQQ,$$

$$V_{\mathbf{F}}PQ =_{df} CPQ.$$

The only primitive rule of procedure is

R1 If P and CPQ then Q .

We shall show how to reduce¹ the number and lengths of the axiom schemes.

It follows at once from a result of Łukasiewicz [3] that *A1-3* may be replaced by the axiom scheme

B1 $CCCPQRCCRPCSP$.

1. The axiom schemes *C* are similar to those obtained by using a general method of Shoesmith [5], but his completeness proof is non-constructive.