

ALGEBRAIC SEMANTICS FOR $S2^0$ AND NECESSITATED EXTENSIONS

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Algebraic techniques are used to show that Feys' system $S2^0$ (cf. [1]) and certain necessitated extensions of $S2^0$, such as Lewis' systems $S2$ and $S3$, have the finite model property, and accordingly are decidable. Representation theorems are then used to establish set-theoretical semantics for the modal systems studied. Where the results obtained are not new they improve on earlier results (such as those of Lemmon in [3]) in two respects; first they provide direct algebraic treatments of the systems, and second they furnish better semantical results (see the discussion of theorem J for $S2$). The techniques used however follow those of McKinsey (in [4]) and Lemmon (in [2] and [3]). Since it is now known that these techniques do not work for all necessitated extensions of $S2^0$, a somewhat piecemeal approach is inevitable. Weak results are also obtained for Feys' system $S1^0$ and Lewis' system $S1$ (for details of these systems see [1]).

The sentential systems studied are of interest not so much as systems containing a viable necessity operator ' \Box ', but as intensional logics which axiomatise epistemic or other operators. For instance $S2^0$ can be interpreted as an epistemic logic such that ' \Box ' reads 'it is believed reasonably that', and $S2$ as an epistemic logic where ' \Box ' reads 'it is known that'. The set-theoretical semantics established are however independent of these epistemic interpretations.

The basic system examined, Feys' $S2^0$, has as postulates:

- T1. $A \ \& \ B \rightarrow A$
- T2. $A \ \& \ B \rightarrow B \ \& \ A$
- T3. $(A \ \& \ B) \ \& \ C \rightarrow A \ \& \ (B \ \& \ C)$
- T4. $A \rightarrow A \ \& \ A$
- T5. $A \rightarrow B \ \& \ B \rightarrow C \rightarrow A \rightarrow C$
- T6. $\Diamond(A \ \& \ B) \rightarrow \Diamond A$

Strict Detachment (SD): $\vdash A, \vdash A \rightarrow B \rightarrow \vdash B$

Adjunction (A): $\vdash A, \vdash B \rightarrow \vdash A \ \& \ B$

Substitutivity of Strict Equivalents (SSE): $\vdash A \leftrightarrow B, \vdash C(A) \rightarrow \vdash C(B)$

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