

## A Model Theoretic Proof of Feferman's Preservation Theorem

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Let  $L$  be a countable first-order language containing a binary relation symbol  $\triangleleft$ . If  $\mathfrak{A}$  and  $\mathfrak{B}$  are  $L$ -structures and  $\mathfrak{A} \subseteq \mathfrak{B}$ , then we say  $\mathfrak{B}$  is a *faithful extension* of  $\mathfrak{A}$  if and only if for any  $a \in \mathfrak{A}$  and  $b \in \mathfrak{B}$  if  $\mathfrak{B} \models b \triangleleft a$ , then  $b \in \mathfrak{A}$ . Thus if  $\triangleleft$  is a linear order on  $\mathfrak{A}$ ,  $\mathfrak{B}$  is a faithful extension if and only if it is an end extension.

In [2] Feferman gives a very natural classification of the formulas which are preserved under faithful extensions. His proof uses a many-sorted interpolation theorem proved by a cut elimination argument. With the introduction of recursively saturated models Barwise and Schlipf [1], and Schlipf [5] attempted to give a unified framework for many preservation and definability theorems. In this note I will give an instructive model theoretic proof of Feferman's theorem. (I should note that Stern [8] and Guichard [4] have given model theoretic proofs of Feferman's theorem using model-theoretic forcing and consistency properties, respectively, but neither of these approaches matches the elegance of [5].)

The proof given here is directly inspired by Friedman's theorem [3] that every countable model of Peano Arithmetic is isomorphic to a proper initial segment of itself and the related embedding results presented in Smoryński [6]. In fact, independently of the author, Smoryński [7] uses Friedman's theorem to prove Feferman's result in the special case that  $\mathfrak{A}$  and  $\mathfrak{B}$  are models of Peano Arithmetic.

### *1 Embedding recursively saturated models*

**Definition 1.1** Let  $L$  be as above. We inductively define  $\Sigma$  a class of  $L$ -formulas as follows:

- (i) If  $\varphi(\bar{v})$  is quantifier free, then  $\varphi(\bar{v})$  is in  $\Sigma$ .
- (ii) If  $\varphi(\bar{v})$  and  $\psi(\bar{v})$  are in  $\Sigma$ , then  $\varphi(\bar{v}) \wedge \psi(\bar{v})$  and  $\varphi(\bar{v}) \vee \psi(\bar{v})$  are in  $\Sigma$ .

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