# Bounds in Weak Truth-Table Reducibility 

KAROL HABART


#### Abstract

A necessary and sufficient condition on a recursive function is given so that arbitrary sets can be truth-table reduced via this function as the bound. A corresponding hierarchy of recursive functions is introduced and some partial results and an open problem are formulated.


Weak truth-table reducibility, often called bounded Turing reducibility, is defined as follows: $A \subseteq \omega$ is weak-truth-table reducible to $B \subseteq \omega\left(A \leq_{\text {wtt }} B\right)$ if there is a recursive function $f$ and an algorithm which answers questions of the form " $n \in A$ ?" when supplied answers to any questions it asks of the form " $m \in B$ ?" for $m \leq f(n)$. The function $f$ is called the bound of the reduction.

The hierarchy of subsets of $\omega$ induced by the relation $\leq_{\text {wtt }}$ was extensively studied in the past (cf. [1]). In this paper, however, a hierarchy of the bounds (i.e. of recursive functions) is considered. We denote by $S(f)$ the set of $A$ such that there is a $B$ such that $A \leq_{\text {wtt }} B$ via a reduction with bound $f$, and we write $f \ll g$ iff $S(f) \subseteq S(g)$. Of course, $\mathcal{R} \subseteq S(f) \subseteq 2^{\omega}$ for all recursive functions $f$. We give necessary and sufficient conditions on $f$ for $S(f)=R$ and for $S(f)=$ $2^{\omega}$, i.e. for $f$ being on the bottom and on the top of the hierarchy induced by $\ll$. We also give a necessary condition for $f \ll g$.

Our interest is focused to the bound of the wtt-reduction by the following phenomenon: A part of an (in general) nonrecursive set $B$ can be given by a list. Having the set $A$ Turing reduced to $B$, a part of $A$ is given which corresponds to the list of a part of $B$, and which may be much larger than the list itself depending mainly on the bound of the reduction.

A motivation for the study of our hierarchy of bounds comes also from the theory of nets of automata. Consider a chain of automata numbered by natural numbers. Suppose each automaton is in one of the two states 0 and 1. Then the state of the whole net is uniquely determined by a set $B \subseteq \omega$ in an obvious way. Now let the automata work, and after some time all of them may stop and the net may come into a state determined by a set $A$. In a fairly devised net we would have $A \leq_{\text {wtt }} B$. The bound $f$ of this reduction depends on how the

