## Bounds in Weak Truth-Table Reducibility

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**Abstract** A necessary and sufficient condition on a recursive function is given so that arbitrary sets can be truth-table reduced via this function as the bound. A corresponding hierarchy of recursive functions is introduced and some partial results and an open problem are formulated.

Weak truth-table reducibility, often called bounded Turing reducibility, is defined as follows:  $A \subseteq \omega$  is weak-truth-table reducible to  $B \subseteq \omega$  ( $A \leq_{wtt} B$ ) if there is a recursive function f and an algorithm which answers questions of the form " $n \in A$ ?" when supplied answers to any questions it asks of the form " $m \in B$ ?" for  $m \leq f(n)$ . The function f is called the bound of the reduction.

The hierarchy of subsets of  $\omega$  induced by the relation  $\leq_{wtt}$  was extensively studied in the past (cf. [1]). In this paper, however, a hierarchy of the *bounds* (i.e. of recursive functions) is considered. We denote by S(f) the set of A such that there is a B such that  $A \leq_{wtt} B$  via a reduction with bound f, and we write  $f \ll g$  iff  $S(f) \subseteq S(g)$ . Of course,  $\Re \subseteq S(f) \subseteq 2^{\omega}$  for all recursive functions f. We give necessary and sufficient conditions on f for  $S(f) = \Re$  and for  $S(f) = 2^{\omega}$ , i.e. for f being on the bottom and on the top of the hierarchy induced by  $\ll$ . We also give a necessary condition for  $f \ll g$ .

Our interest is focused to the bound of the wtt-reduction by the following phenomenon: A part of an (in general) nonrecursive set B can be given by a list. Having the set A Turing reduced to B, a part of A is given which corresponds to the list of a part of B, and which may be much larger than the list itself—depending mainly on the bound of the reduction.

A motivation for the study of our hierarchy of bounds comes also from the theory of nets of automata. Consider a chain of automata numbered by natural numbers. Suppose each automaton is in one of the two states 0 and 1. Then the state of the whole net is uniquely determined by a set  $B \subseteq \omega$  in an obvious way. Now let the automata work, and after some time all of them may stop and the net may come into a state determined by a set A. In a fairly devised net we would have  $A \leq_{wtt} B$ . The bound f of this reduction depends on how the

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