

## Preservation by Homomorphisms and Infinitary Languages

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**Abstract** In this paper we study when sentences of infinitary languages are preserved by homomorphisms. This is done by using generalized Henkin construction. By the same technique we can also study when a sentence has an equivalent sentence which is in normal form.

The so-called Hintikka game, which is a straightforward generalization of Henkin construction, is used in, e.g., Hyttinen [1], [2] and Oikkonen [6]. In this paper we refine the game and prove a preservation theorem by using it.

Throughout this paper we assume that  $\kappa$  is weakly compact. This is done because in the proofs we construct certain trees, which have no branches of length  $\geq \kappa$  and no nodes that have  $\geq \kappa$  immediate successors, and then under the assumption that  $\kappa$  is weakly compact we know that the trees are of cardinality  $< \kappa$ .

We begin this chapter by defining the language  $M_{\kappa\kappa}$ . This language was first defined and studied by M. Karttunen in [4]. By a  $\lambda$   $\kappa$ -tree  $T$  we mean a tree such that each node has  $< \lambda$  immediate successors, there are no branches of length  $\geq \kappa$  and if  $x$  and  $y$  are limit nodes and  $\{t \in T \mid t < x\} = \{t \in T \mid t < y\}$  then  $x = y$ .

**1.1 Definition**  $\phi = (T, l)$  is a formula of  $M_{\kappa\kappa}$  if

- (1)  $T$  is a  $\kappa\kappa$ -tree without branches of limit length, i.e. every branch has a maximal element;
- (2)  $l$  is a labeling function with the properties
  - a: if  $t \in T$  does not have any successors then  $l(t)$  is either an atomic or negated atomic formula;
  - b: if  $t \in T$  has exactly one immediate successor then  $l(t)$  is of the form  $\exists x$  or  $\forall x$ ,  $x$  variable;
  - c: if  $t \in T$  has more than one immediate successor then  $l(t)$  is either  $\vee$  or  $\wedge$ .

*Received October 30, 1989; revised January 15, 1990*