

## Enumerations of Turing Ideals with Applications

DAVID MARKER\*

**Abstract** We examine enumerations of ideals in the Turing degrees and give several applications to the model theory of first- and second-order arithmetic.

A *Turing ideal* is a collection of subsets of  $\omega$  closed under Turing reducibility and join. If  $I$  is a countable Turing ideal we say that  $E$  is an *enumeration* of  $I$  if and only if  $I = \{E_n : n \in \omega\}$  where  $E_n = \{m : \langle n, m \rangle \in E\}$ . Enumerations of Turing ideals play an important role in the study of degrees coding recursively saturated models of Peano Arithmetic (see [5]). Our goal in this paper is to point out some simple facts about enumerations of Turing ideals and examine their consequences in the model theory of first- and second-order arithmetic.

**1  $\omega$ -incompleteness theorems** We consider three subsystems of second-order arithmetic,  $RCA_0$ ,  $ACA_0$ , and  $WKL_0$ .  $RCA_0$  is axiomatized by  $P^-$  (Peano Arithmetic without the induction axioms), the axiom of extensionality and the schemas of  $\Sigma_1^0$ -induction and recursive comprehension.  $WKL_0$  is obtained from  $RCA_0$  by adding an axiom saying every infinite subtree of  $2^{<\omega}$  has an infinite path.  $ACA_0$  is obtained from  $RCA_0$  by adding the schema of arithmetic comprehension. For further information on these theories the reader should consult [7].

Our recent interest in this subject was motivated by considering the following incompleteness theorem of Steel.

**Theorem 1.1** (Steel [9]) *Let  $T$  be an  $\omega$ -consistent arithmetic extension of  $ACA_0$ . There is an  $\omega$ -model  $\mathbf{M}$  of  $T$  such that  $\mathbf{M} \models$  “there is no  $\omega$ -model of  $T$ ”.*

Thus even in  $\omega$ -logic  $T$  does not prove its own  $\omega$ -consistency. Steel’s result is actually much stronger. Suppose  $\mathbf{M} = (\omega, X)$  and  $\mathbf{N} = (\omega, Y)$  are models of  $RCA_0$ . We say that  $\mathbf{M} \gg \mathbf{N}$  if and only if there is  $E \in X$  an enumeration of  $Y$ . If  $\mathbf{M} \models$  “there is an  $\omega$ -model of  $T$ ”, then there is an  $E \in X$  such that  $\mathbf{N} =$

---

\*Partially supported by an NSF grant.