

FINITE LIMITATIONS ON DUMMETT'S LC

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The propositional system **LC** of [1] can be based on axioms for \supset (implication), \wedge (conjunction), a constant **f**, and definitions for \vee (alternation) and \neg (negation), as hereunder. In primitive notation, elementary variables and **f** are wffs, and if α , β are wffs so are $(\alpha \supset \beta)$, $(\alpha \wedge \beta)$. To restore primitive notation in the sequel, replace dots by left parentheses with right terminal mates; in a sequence of wffs separated only by implications, restore parentheses by left association; enclose the whole in parentheses. If S is a system, S_c is its implicational fragment, containing only variables and implications. If α is provable (not provable) in S , we write $\frac{}{S} \alpha$ ($\frac{}{S} \neg \alpha$); if α is uniformly valued 0 (is not uniformly valued 0) by the matrix \mathfrak{M} , we write $\frac{}{\mathfrak{M}} \alpha$ ($\frac{}{\mathfrak{M}} \neg \alpha$). As a basis for **LC** we take, with detachment and substitution, the axioms and definitions:

- 1 $p \supset . q \supset p$
- 2 $p \supset (q \supset r) \supset . p \supset q \supset . p \supset r$
- 3 $p \supset q \supset r \supset . q \supset p \supset r \supset r$
- 4 $\mathbf{f} \supset p$
- 5 $(p \wedge q) \supset p$
- 6 $(p \wedge q) \supset q$
- 7 $p \supset . q \supset (p \wedge q)$

Def. \vee $(\alpha \vee \beta) = (\alpha \supset \beta \supset \beta) \wedge (\beta \supset \alpha \supset \alpha)$

Def. \neg $\neg \alpha = \alpha \supset \mathbf{f}$

[2] shows that 1-3 suffice for \mathbf{LC}_c , and it is well known that 1-2 suffice for \mathbf{IC}_c , the positive logic. By [1] the infinite adequate matrix for **LC** is $\mathfrak{M} = \langle M, \{0\}, \wedge, \supset, \mathbf{f} \rangle$ where $M = \{0, 1, 2, \dots, \omega\}$ and

$$\begin{aligned} a \wedge b &= \max(a, b), \\ a \supset b &= \begin{cases} 0 & \text{if } a \geq b, \\ b & \text{if } a < b, \end{cases} \\ \mathbf{f} &= \omega. \end{aligned}$$