

ON A RECENT ALLOTMENT OF PROBABILITIES TO
OPEN AND CLOSED SENTENCES

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Probabilities, though frequently allotted to closed sentences, have rarely been allotted to open ones. The recent scheme by Kemeny, Mirkil, Snell, and Thompson in *Finite Mathematical Structures* for allotting probabilities to sentences of the form ' $f(x) = a$ ' is therefore of considerable interest.¹ It has, however, a shortcoming which I should like to discuss here and, possibly, remedy.

Let ' f ' be a functional constant, ' x ' an individual variable, and ' a ' an individual constant; let U be the (finite) set of values of ' x '; and let A be the subset of U whose members satisfy ' $f(x) = a$ '. Kemeny *et al.* then take the probability of ' $f(x) = a$ ' to be $m(A)$, where $m(A)$ is the measure (in some appropriate sense of the word 'measure') of A .² Their scheme is attractive enough and mirrors to some extent what mathematicians understand by the probability of a set.³ Kemeny *et al.* are careful, of course, to restrict it to open sentences of the form ' $f(x) = a$ '. Consider, however, a closed sentence of the kindred form ' $f(b) = a$ ', where ' b ' is an individual constant. Since ' $f(b) = a$ ' does not contain any occurrence of ' x ', it would normally be held to be satisfied by every member of U when true, by none when false. One would accordingly expect Kemeny *et al.* to take the probability of ' $f(b) = a$ ' to be 1 when ' $f(b) = a$ ' is true, 0 when ' $f(b) = a$ ' is false. Yet in their scheme for allotting probabilities to closed sentences, a scheme I shall go into below, they let the probability of a closed sentence equal 1 only when the sentence is logically true, 0 only when it is logically false.⁴ ' $f(b) = a$ ' being neither logically true nor logically false, its probability must therefore differ by that scheme from either one of 1 and 0, a disturbing enough result.

The difficulty becomes even more acute when the calculus, call it C , to whose sentences probabilities are allotted is a simple applied predicate calculus of the first order with identity.

Assume indeed that a set D of individuals has been singled out as the domain of C , a member of D paired with each individual constant W of C as the individual designated by W , and a class of ordered n -tuples of members of D paired with each n -adic predicate constant F of C as the extension of