

GENTZEN-LIKE SYSTEMS FOR PARTIAL
 PROPOSITIONAL CALCULI: I

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1. *Introduction.* The existence of many decidable partial propositional calculi leads to a desire for the construction of Gentzen-like systems (with provable cut-elimination theorems) for such calculi. Two problems arise in the construction of such systems. First, one must to some extent formalize what one means by a Gentzen-like system. Second, one must find appropriate deduction rules for the systems so that the cut-elimination theorem can be proved. Here we consider the general definition and the two simpler one axiom sub-partial propositional calculi of the system of [1]. Part II of this paper will consider the other one axiom subcalculus.

2. *Classical Gentzen-like Systems.* The classical Gentzen-like systems, for example [2], might be characterized as follows.

- (i) The objects manipulated are n -tuples of well-formed formulae.
- (ii) Certain of these n -tuples with (easily) effectively recognizable forms are taken as axioms.
- (iii) Deduction rules each take some n -tuple (or pair of n -tuples) and make (easily) effectively recognizable changes in the n -tuples. Most of these changes are permutations of the symbols of the n -tuples or the introduction of a new symbol or formula into the n -tuple.
- (iv) A certain (easily) effectively recognizable subset of the n -tuples constructed from the axioms by the deduction rules is identifiable with the set of theorems of the classical theory for which the Gentzen-like system is being constructed.
- (v) Often a theorem (cut-elimination) is provable on the basis of which one can effectively construct, given a n -tuple of the proper form for (iv), a finite set of possible proofs for it in such a way that one can actually decide if it is a theorem by determining if one of these possible proofs is actually a proof.

3. *Definitions.* Given a partial propositional calculus C with language L with the usual variables, constants, connectives and formation rules, let #