

TERMINAL FUNCTORS PERMISSIBLE
 WITH SYLLOGISTIC

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(Notes of a paper delivered in July 1957, on the possibilities of enriching syllogistic as axiomatised, e.g. in Łukasiewicz's *Aristotle's Syllogistic*, with complex terms, with appropriate axioms).

I. *General conclusions**

(i) The *logical product* of a pair of terms, kab , does not lend itself for use in this way, since one would expect it to have $\vdash Akaba$ and $\vdash Akabb$, which would yield $\vdash Iab$ by Darapti, making all I propositions true.

(ii) Similarly with the *null term* \wedge ; one would expect this to have $\vdash A \wedge a$, but $\vdash A \wedge a$ and $\vdash A \wedge b$ (obtained from $\vdash A \wedge a$ by substitution) would yield Iab , again by Darapti.

(iii) *Disjunction* and *implication* are possible, so long as they are not combined with negation (when of course conjunction and the null term would be definable).

(iv) The *modal* λ and μ ("necessarily", "possibly") seem alone to combine with negation, n .

II. *Postulates for non-modal terminal functors*

(i) If we take the syllogistic E as undefined (defining Aab as $Eanb$ and I as NE), syllogistic with terminal negation may be axiomatised by 1. $Eana$, 2. $NEaa$ and 3. $Camenes$.

(ii) For terminal disjunction, v , we may subjoin to ordinary syllogistic (without n) first of all the four axioms

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| 1. $Aavab$ | 3. $CAacCAbcAvabc$ |
| 2. $Abvab$ | 4. $CEacCEbcEvabc$ |

Where t and s are terms, we write $t = s$ where we have both $\vdash Ats$ and $\vdash Ast$; and where p and q are propositions, we write $p = q$ where we have both $\vdash Cpq$ and $\vdash Cqp$. The above for v will give $vaa = a$, $Avabc = KAacAbc$, and $Evabc = KEacEbc = Ecvab$. The addition of

*See Historical Note by A. N. Prior added to this paper.