

THE CHURCH ROSSER THEOREM FOR STRONG REDUCTION
 IN COMBINATORY LOGIC

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The Church Rosser theorem concerns a property relating to certain preordering relations [2a]. Originally it was stated for lambda conversions in a paper by Church and Rosser [1].

To define the property let Γ be a preordering and $=$ be its symmetric closure. The property in question states

(C R). *If $M = N$, then there is an L such that $M \Gamma L$ and $N \Gamma L$.*

In this paper we give a proof that strong reduction (as modified by the author in a previous paper [3]) has the property (C R). For strong reduction, the symmetric closure is simply combinatory logic with equality [2b]. The following results were proved in [3] and [5] and will be used here.

Lemma 1. If $X = [x]x$, then $\lambda x.Xx \succ \lambda x.x$ by Type I steps only. In other words, the contractum of a Type III step may be reversed to the original redex by a single Type II step followed by Type I steps.

Lemma 2. The contraction of a Type II redex P may be reversed provided there are no intervening steps interior to the contractum of P .

Lemma 3. (Theorem 2.II of [5]) If there is a standard reduction from M_0 to M_n and if there is a single step of Type I or III from M_0 to N_0 , then there is an N_n such that there is a standard reduction $N_0 \succ N_n$ and M_n .

Lemma 4. (Lemma 5 of [5]) If there is a reduction from M_0N_0 beginning with a Type II step yielding $(\lambda x.M_1)N_0$ and continuing to $(\lambda x.M_m)N_n$, then there is a reduction from M_0N_0 to $[N_n/x]M_m$ (where $[N_n/x]M_m$ means the substitution of N_n for x in M_m) in exactly the same number of steps.

Lemma 5. (Theorem 3 of [5]) If there is a strong reduction from X to Y where neither X nor Y contain lambda expressions, then there is a Z such that there is a standard reduction from X to Z and $Y \succ Z$.

Lemma 6. (Corollary A of [5]) If there is a strong reduction from M to N , then there is a Z such that there is a reduction consisting of zero or