

NORMAL AND SKEW SYSTEMS

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Preface Suppose the formal system T has a binary connective \rightarrow . If T has:

$$\mathbf{MP}: \frac{A, A \rightarrow B}{B}$$

among its rules of inference, it is natural to ask: "Which theorems of T are still provable (without repetitions) if one refuses to use **MP** except when A precedes $A \rightarrow B$?" Roughly speaking, we will say that T is normal if the answer to this question is "All of them."

In forming a precise definition a certain difficulty becomes apparent. One might be able to augment the sequence $\dots, A \rightarrow B, \dots, A, \dots, B$ by new formulas to obtain $\dots, A \rightarrow B, \dots, A, \dots, C, \dots, C \rightarrow B, \dots, B$, and thus avoid the issue. For example, if T had the axiom schemes:

- A1) $A \rightarrow (B \rightarrow A)$
 A2) $(A \rightarrow (B \rightarrow C)) \rightarrow ((A \rightarrow B) \rightarrow (A \rightarrow C)),$

then $\dots, A \rightarrow B, \dots, A, \dots, B$ could be replaced by:

⋮

$A \rightarrow B$

⋮

A

⋮

F

$A \rightarrow (F \rightarrow A)$

$F \rightarrow A$

$(A \rightarrow B) \rightarrow (F \rightarrow (A \rightarrow B))$

$F \rightarrow (A \rightarrow B)$

$(F \rightarrow (A \rightarrow B)) \rightarrow ((F \rightarrow A) \rightarrow (F \rightarrow B))$

$(F \rightarrow A) \rightarrow (F \rightarrow B)$

$F \rightarrow B$

A1

MP

A1

MP

A2

MP

MP