

A FORMAL SYSTEM FOR THE NON-THEOREMS  
 OF THE PROPOSITIONAL CALCULUS

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*Introduction* The completeness of the classical propositional calculus allows us to give a deductive system consisting of finitely many *axiom schemas* and finitely many *rules of inference*, that permit us to pass from a formula or a pair of formulae to a syntactically related formula, in such a manner that the formulae obtained inductively from the axioms by repeated application of the rules are exactly the tautologies. In this paper we give an analogous deductive system (more concretely, a Hilbert type system) such that the formulae deduced are exactly those that *are not* tautologies, the non-theorems of the propositional calculus. Obviously, this has to be the most non-standard of the non-classical logics. It is important to note that there are many other algorithms to generate recursively the non-theorems, since the propositional calculus is decidable. Usually they are based in the methodical search for a counterexample, but they lack the inductive character of a Hilbert type system, where every formula involved in a deduction is itself deducible. In our system, unlike semantic tableaux or refutation trees, every formula introduced in a deduction is a non-tautology, and it is introduced only if it is a non-tautological axiom, or it follows by one of the non-tautological rules of inference from non-tautologies introduced earlier in the deduction.

1 *Axioms and rules* We assume that the only connectives are  $\sim$  and  $\supset$ .  $p, q, p_1, p_2, \dots$  denote atomic formulae.  $\alpha, \beta, \gamma, \dots$  denote arbitrary formulae. We define  $\mathcal{P}(\alpha) = \{p \mid p \text{ occurs in } \alpha\}$ .

Axioms

- A1  $p \supset \sim p$  ( $p$  atomic)  
 A2  $\sim p \supset p$  ( $p$  atomic)

Rules

- R1 (a)  $\frac{\alpha}{p \supset \alpha}$  ( $p$  atomic,  $p$  does not occur in  $\alpha$ )

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