

NEGATION, MATERIAL EQUIVALENCE, AND CONDITIONED  
NONCONJUNCTION: COMPLETENESS AND DUALITY

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1 *Object of paper* The object of this paper is threefold: to prove the functional incompleteness of  $\{\sim, \equiv\}$  without appeal to a tedious analysis of cases; to give a proof simpler than the one in [2], p. 284 f., of the non-existence of indigenous Sheffer connectives for  $\{\sim, \equiv\}$ ; and to furnish self-dual (in Church's sense) ternary Sheffer connectives for propositional logic.

2 *Functional incompleteness of  $\{\sim, \equiv\}$* 

Lemma 1 *Let  $A$  be a propositional wff containing no connectives other than  $\sim$  and  $\equiv$ , and let  $A'$  be the wff that results when all occurrences of  $\sim$  in  $A$  are deleted. Then  $A$  is equivalent to  $A'$  or to  $\sim A'$ .*

*Proof:* Let  $\Leftrightarrow$  signify (semantic) equivalence. Since  $B \Leftrightarrow \sim \sim B$  and since  $\sim(B \equiv C) \Leftrightarrow \sim B \equiv C \Leftrightarrow B \equiv \sim C$ , Lemma 1 follows from the substitutivity of equivalents by induction.

Theorem 1  $\{\sim, \equiv\}$  is functionally incomplete.

*Proof:* We call a truth-value  $\alpha$  a *fixed point* for a wff  $B$  just in case the value of  $B$  is  $\alpha$  when all its variables are assigned the value  $\alpha$ , and we say that  $\alpha$  is a *fixed point* for an  $n$ -ary connective  $\otimes$  just in case  $\alpha$  is a fixed point for  $\otimes(p_1, \dots, p_n)$ . Notice that **t** is a fixed point for  $\equiv$  and for any wff that contains no connectives other than  $\equiv$ . Suppose that  $\&$  is definable from  $\{\sim, \equiv\}$ . Then by Lemma 1 there is a wff  $A(p, q)$  containing no connectives other than  $\equiv$  such that  $p \& q$  is equivalent to  $A(p, q)$  or to  $\sim A(p, q)$ . But since **t** is a fixed point for both  $p \& q$  and  $A(p, q)$ ,  $p \& q$  cannot be equivalent to  $\sim A(p, q)$ . So,  $p \& q \Leftrightarrow A(p, q)$ . Therefore,  $p \& \sim q \Leftrightarrow A(p, \sim q)$ . By Lemma 1 again,  $A(p, \sim q)$  is equivalent to  $A(p, q)$  or to  $\sim A(p, q)$ . Since **t** is not a fixed point for  $p \& \sim q$ , we have  $A(p, \sim q) \Leftrightarrow p \& \sim q \Leftrightarrow \sim A(p, q)$ . So,  $\sim(p \& \sim q) \Leftrightarrow A(p, q) \Leftrightarrow p \& q$ , which is a contradiction. So  $\&$  is not definable from  $\{\sim, \equiv\}$ . Theorem 1 follows.

3 *Indigenous Sheffer connectives* An  $n$ -ary connective  $\otimes$  is said to be a