

GENERAL COMPUTABILITY

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1 Introduction A theory of computability consists of a domain A of cardinality greater than one and a definition of computability for multivariable partial functions from A into A . The class \mathcal{F} of all partial functions satisfying that definition is called the class of computable functions. In this paper, we consider some properties of \mathcal{F} which have been of importance in various specific theories. In particular, we consider conjunctions of the following properties:

- (1) \mathcal{F} is the closure of a family \mathcal{G} under composition.
- (2) \mathcal{F} is closed under piecewise composition.
- (3) \mathcal{F} is the closure of a family \mathcal{G} under composition and iteration where the iterate f^∞ of a partial function f is $\lambda x [f^n(x)]$, where n denotes $\mu m [f^m(x) = f^{m+1}(x)]$
- (4) \mathcal{G} is closed under piecewise composition.
- (5) \mathcal{G} is finite.
- (6) \mathcal{F} contains the test function $\lambda w, x, y, z [x \text{ if } w = z; \text{ else } y]$.
- (7) \mathcal{F} includes a pairing system, i.e., a set $\{\rho, \sigma, p\}$ of partial functions such that $\rho(p(a, b)) = a$ and $\sigma(p(a, b)) = b$ for all a and b in A .
- (8) ρ, σ , and the members of \mathcal{G} are total functions on A .
- (9) \mathcal{F} has an indexing; i.e., \mathcal{F} is closed under composition and contains a partial function $*$ such that every member of \mathcal{F} is of the form:

$$\lambda x_1, \dots, x_n [(\dots (a * x_1) * \dots * x_n)]$$

for some a in A .

- (10) \mathcal{F} has a uniform indexing in the sense that \mathcal{F} has an indexing $*$ such that every member of \mathcal{F} is of the form

$$\lambda x_1, \dots, x_n [g(x_1, \dots, x_{n-1}) * x_n]$$

for some total g in \mathcal{F} .