

NOTE ON A STRONG LIBERATED MODAL LOGIC AND ITS
RELEVANCE TO POSSIBLE WORLD SKEPTICISM

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We consider a modal language with a set of sentence parameters $P = \{P_1, P_2, \dots\}$ and the symbols L (necessity), \supset , and \neg for sentential connectives. We will consider axiomatic systems based on the following axiom schemes:

A1 All tautologies

A2 $L(E \supset E') \supset (LE \supset LE')$.

As rules of proof, we assume the following:

MP If E and $E \supset E'$ are provable, then so is E' .Nec If E is provable, then so is LE .

Consider the two following axiom schemes:

A3 $LE \supset E$ A4 $E \supset LE$.

If we add both A3 and A4 to our basic system, we obtain the trivial system in which for any expression E , LE is equivalent to E . The classical systems adopt A3 and exclude A4 to obtain **T** and stronger systems (see [1]). Murungi [6] has given a formal proof of the claim that the system obtained by adding A4 to the basic system, sketched above, is consistent. The system with A4 is called **T'**.

While it is true that **T'** is consistent and does not collapse into the trivial system, **T'** comes *very* close to collapsing into the trivial system, as we shall show by providing a semantics and completeness proof for **T'**. The system **T'** turns out to be one of the so-called liberated modal logics. It is a very strong logic containing all of the liberated systems discussed in [2]-[5].

The semantic range is the usual $\{t, f\}$. We designate by G an indexed set of functions g mapping the set of sentence parameters P into the semantic range. An interpretation is an ordered triple (g, W, R) where: $g \in G$; $W \subseteq G$; R is a binary relation (accessibility) defined on W ; if W is not empty, then