

## A Note on the Hanf Number of Second-Order Logic

MATT KAUFMANN\*

The Hanf number of a logic  $L$  is the least cardinal  $\kappa$  such that every sentence of  $L$  that has a model of power at least  $\kappa$  has arbitrarily large models. The Hanf number  $\kappa^\Pi$  of second-order logic is very large; for example, it is readily seen to exceed the first measurable cardinal (if there is one). In fact, Barwise [1] showed that one cannot prove that  $\kappa^\Pi$  exists within the theory  $ZF_1$ , which is  $ZFC$  with the full subset schema but with collection only for  $\Sigma_1(\mathcal{P})$  formulas. (Here  $\mathcal{P}$  is a unary function symbol, and the power set axiom reads:  $\forall x \forall y (y \subseteq x \leftrightarrow y \in \mathcal{P}(x))$ .) Moreover, within  $ZF_1$  he proved that  $R_{\kappa^\Pi} \models ZF_1$ , and in fact  $\kappa^\Pi$  is the  $(\kappa^\Pi)^{\text{th}}$  cardinal with this property. Friedman [3] improved this result by showing that even in the weaker theory  $T = ZF_0 + (\beta)$ , where  $ZF_0 = KP(\mathcal{P}) + [\text{Power set axiom}]$ , if  $\kappa^\Pi$  exists then  $R_{\kappa^\Pi} \prec_{\Sigma_1(\mathcal{P})} V$ .<sup>1</sup>

In this short note we use Friedman's result to give a new characterization of  $\kappa^\Pi$  (Theorem 1 below). A related characterization is given in Väänänen [5] (Corollary 5.7):

$$(1) \quad \kappa^\Pi = \sup\{\alpha : \alpha \text{ is } \Sigma_2\text{-definable}\},$$

where a set  $S$  is  $\Sigma_2$ -definable if the predicate " $x \in S$ " is a  $\Sigma_2$ -definable predicate of set theory.<sup>2</sup> (Väänänen's result is actually more general.) Here is an outline of a proof of (1). For  $\geq$ , if  $\phi(x)$  is a  $\Sigma_2$  (or  $\Sigma_1(\mathcal{P})$ ) definition of " $x \in \alpha$ " then consider the following sentence  $\psi$  of second-order logic, which holds in  $(R_\kappa, \in)$  if  $\kappa$  is least such that  $\phi(x)$  defines " $x \in \alpha$ " in  $(R_\kappa, \in)$ :

$$\psi \equiv \text{"The universe is of the form } (R_\delta, \in)\text{"} \wedge (\exists! \beta)(\forall x)(\phi(x) \leftrightarrow x \in \beta) \\
\wedge \forall \beta [\forall x(\phi(x) \leftrightarrow x \in \beta) \rightarrow \forall \gamma \exists x \in \beta (R_\gamma \models \neg \phi(x))].$$

Then  $\psi$  has a model of power at least  $|\alpha|$ , but it's easy to see that  $\psi$  does not have arbitrarily large models. For  $\leq$ , observe that if  $\phi$  is a sentence of second-

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