## The Simple Consistency of a Set Theory Based on the Logic CSQ

## **ROSS T. BRADY\***

This paper proves the simple consistency of the set theory *CST*. *CST* has the Generalized Comprehension Axiom (*GCA*),  $(\exists y)(\forall x)(x \in y \leftrightarrow A)$ , and the Extensionality Rule,  $x = y \Rightarrow x \in w \leftrightarrow y \in w$ , where  $x = y =_{df} (\forall z)(z \in x \leftrightarrow z \in y)$ . *CST* is based on a logic *CSQ*, which is semantically described below.

## CSQ Primitives

- 1.  $\sim, \&, \rightarrow, \forall$  (connectives and quantifier)
- 2.  $f, g, h, f', \ldots$  (predicate constants)
- 3.  $x, y, z, x', \ldots$  (individual variables)
- 4.  $a_1, a_2, a_3, a_4, \ldots$  (individual constants).

## CSQ Formulas

- 1. An individual variable or constant is a term.
- 2. If  $t_1, \ldots, t_n$  are terms and f is a predicate constant, then  $ft_1 \ldots t_n$  is an atomic formula.
- 3. If A and B are formulas and x is an individual variable then  $\sim A$ , A & B,  $A \rightarrow B$  and  $(\forall x)A$  are formulas.

A sentence is a formula with no free variables.

A CSQ model structure (CSQ m.s.) consists of ordered triples  $\langle T, K, R \rangle$ , such that K is a set, T is a member of K, and R is a two-place relation on K, with the following postulates holding: For  $\alpha \in K$ ,

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