

Post's Functional Completeness Theorem

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Abstract The paper provides a new proof, in a style accessible to modern logicians and teachers of elementary logic, of Post's Functional Completeness Theorem. Post's Theorem states the necessary and sufficient conditions for an arbitrary set of (2-valued) truth functional connectives to be expressively complete, that is, to be able to express every (2-valued) truth function or truth table. The theorem is stated in terms of five properties that an arbitrary connective may have, and claims that a set of connectives is expressively complete iff for each of the five properties there is a connective that lacks that property.

Everyone knows the technique whereby, given an arbitrary (2-valued) truth table, one can construct a conjunctive (or disjunctive) normal form formula (using only connectives from $\{\vee, \wedge, \sim\}$) which has exactly that truth table. This proves only that the set of connectives $\{\vee, \wedge, \sim\}$ is *functionally complete*: any (2-valued) truth table can be constructed from them. Everyone also knows the definitions of \wedge in terms of $\{\vee, \sim\}$ and of \vee in terms of $\{\wedge, \sim\}$. This shows that $\{\wedge, \sim\}$ and $\{\vee, \sim\}$ are also functionally complete sets of connectives. Everyone also knows that the sheffer stroke functions, \uparrow and \downarrow , are each functionally complete. Most everyone knows that $\{\rightarrow, \mathbf{F}\}$ is functionally complete and that $\{\rightarrow, \underline{\vee}\}$ is functionally complete (\mathbf{F} is the constant-false truth function, $\underline{\vee}$ is "exclusive or"). Some people, having worked through Church ([1], p. 131f.), even know that $\{\{\}, \mathbf{T}, \mathbf{F}\}$ is functionally complete ($\{\}$ is the ternary connective of "conditional disjunction": $[p, q, r]$ means "if q , then p else r "). However, what is not generally known is why these things are so. What is it about these particular sets of connectives that makes them functionally complete while (say) $\{\leftrightarrow, \sim\}$ is not functionally complete?

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