

## COUNTING FINITE SUBSETS OF AN IMMUNE SET

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ABSTRACT. Let  $P$  be a recursive property of finite sets (of nonnegative integers) and  $\nu$  an immune set of RET (i.e., recursive equivalence type)  $N$ . Consider the question, "How many finite subsets of  $\nu$  have property  $P$ ?" We shall answer this question if  $P$  has the additional property that  $P(\alpha)$  if and only if  $P(\beta)$ , for every two finite sets  $\alpha$  and  $\beta$  of the same cardinality.

**1. Preliminaries.** We use the word *number* for nonnegative integer, *set* for collection of numbers and *class* for collection of sets. The set of all numbers and the empty set are denoted by  $\varepsilon$  and  $o$  respectively.  $V$  stands for the class of all sets,  $Q$  for the class of all finite sets,  $\subset$  for inclusion and  $\subset_+$  for proper inclusion. If  $f$  is a function,  $\delta f$  and  $\rho f$  denote its domain and range respectively. Also,  $f_n$  means the same as  $f(n)$ . The cardinality of a collection  $\Gamma$  is denoted by  $\text{card}(\Gamma)$  or  $\text{card } \Gamma$ . The reader is assumed to be familiar with the basic properties of the collection  $\Lambda$  of all isols. For a survey of basic results see §1 of [2]; for a detailed exposition, see [4] or [5]. We write  $\alpha \sim \beta$  for  $\alpha$  is equivalent to  $\beta$ , i.e.,  $\text{card } \alpha = \text{card } \beta$  and  $\alpha \simeq \beta$  for  $\alpha$  is recursively equivalent to  $\beta$ , i.e.,  $\text{Req}(\alpha) = \text{Req}(\beta)$ . Some properties of combinatorial operators will be used; these are discussed in [1], [4] and [5]. The class  $Q$  of all finite sets can be effectively generated without repetitions in an infinite sequence. We shall use a particular sequence of this type, the so-called *canonical enumeration*  $\langle \rho_n \rangle$ ; see [2, p. 277]. For each finite set  $\sigma$  there is exactly one number  $i$  such that  $\sigma = \rho_i$ ; this number is called the *canonical index* of  $\sigma$  and is denoted by  $\text{can}(\sigma)$  or  $\text{can } \sigma$ . If  $S \subset Q$  we write  $\text{can } S$  for  $\{\text{can } \sigma \mid \sigma \in S\}$ . The function  $r_n = \text{card } \rho_n$  is recursive. For  $n, k \in \varepsilon$  and  $\alpha \in V$ ,

$$\begin{aligned} \nu_n &= \{x \in \varepsilon \mid x < n\}, & 2^\alpha &= \{x \in \varepsilon \mid \rho_x \subset \alpha\}, \\ [\alpha; k] &= \{x \in 2^\alpha \mid r_x = k\}, & \int [n; k] &= \text{card } [\nu_n; k]. \end{aligned}$$

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