

A NOTE ABOUT CONNECTION OF THE FIRST-ORDER  
FUNCTIONAL CALCULUS WITH MANY VALUED  
PROPOSITIONAL CALCULI

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By virtue of a generalization of the satisfiability definition, see [2], we described in [3] an approximation of the first-order functional calculus by Boolean many valued propositional calculi in which the quantifier  $\Pi$  had a finite meaning.

In this paper we shall describe another approximation of the calculus by many valued Boolean propositional calculi based in [4]; the proof of the approximation is analogical to [3] and it is given in [5].

We consider here a Boolean algebra with operations  $\neg$ /complementation/,  $+$  /addition/ and with elements which are  $n$ -tuples  $(w_1, \dots, w_n)$  of numbers 0 and 1.

We use notations of [3] and especially the following:

1. variables of the calculus:
  - (1') free:  $x_1, \dots$  /simply  $x$ /,
  - (2') apparent:  $a_1, \dots$  /simply  $a$ /.
2. relations signs:  $f_1, \dots, f_c$ ;  $\bar{c}$  - maximum of arguments of ones.
3.  $w(E)$  - the number of different free / $p(E)$  - apparent/variables occurring in  $E$ .
4.  $i(E)$  - maximum of indices of those and only those variables which occur in  $E$ .
5.  $n(E) = i(E) + p(E)$ .
6.  $E(u/z)$  - substitution of  $u$  for each occurrence of  $z$  in  $E$  /with knowing conditions/.
7.  $C\{E\}$  - the set of all significant parts of  $E$ :  
 $H \in C\{E\} \equiv H = E$  or there exist  $E_1 \in C\{E\}$ ,  $F, G, H_1$  such that:  
 $(H = F) \wedge (E_1 = F') \vee \{(H = F) \vee (H = G)\} \wedge (E_1 = F + C) \vee (\exists i)\{H = H_1(x_i/a)\} \wedge (E_1 = \Pi aH_1)$ .
8.  $S(k)$  - the set of all atomic formulas  $R$  such that indices of free variables occurring in  $R$  are  $\leq k$ .
9.  $Q$ -function on  $S(k)$  with values  $n$ -tuples  $(w_1, \dots, w_n)$  of numbers 0 and 1.

**D.1.**  $g(k, j, t, Q, m) \equiv (k \leq m) \wedge (R) \{ (R \in S(k)) \wedge \{ Q(R) = (w_1, \dots, w_n) \text{ for some } w_1, \dots, w_n \} \rightarrow (w_j = w_t) \}$

By means of the function  $Q$  we give an inductive definition of the functional  $V$  which is defined for an arbitrary formula  $E$  such that  $i(E) \leq k$  and  $k + p(E) \leq m$ :

- (1d)  $V\{k, Q, m, R\} = Q(R)$ , if  $R \in S(m)$ ,
- (2d)  $V\{k, Q, m, F'\} = V^{-1}\{k, Q, m, F\}$ ,
- (3d)  $V\{k, Q, m, F + G\} = V\{k, Q, m, F\} \dot{+} V\{k, Q, m, G\}$ ,
- (4d)  $V\{k, Q, m, \Pi aF\} = (w_1, \dots, w_n)$ , for some  $w_1, \dots, w_n \equiv \dots$   
 $(j)\{(j \leq n) \rightarrow (w_j = 1) \equiv \dots$   
 $(r)\{(r \leq k) \wedge (V\{k, Q, m, F(x_r/a)\} = (w_1^r, \dots, w_n^r) \text{ for some } w_1^r, \dots, w_n^r) \rightarrow (w_j^r = 1)\} \wedge (t)\{(t \leq n) \wedge g(k, j, t, Q, m) \wedge \{V\{k+1, Q, m, F(x_{k+1}/a)\} = (v_1, \dots, v_n) \text{ for some } v_1, \dots, v_n\} \rightarrow (v_t = 1)\} \} .$

**D.2.**  $J(Q, m, G) \equiv (k)\{(i(G) \leq k) \wedge (k + p(G) < m) \rightarrow (V\{k+1, Q, m, G\} \subset V\{k, Q, m, G\})\}$ .

**D.3.**  $F \in P(Q, m, E) \equiv (\exists G)\{(G \in C\{E\}) \wedge (J(Q, m, G) \rightarrow V\{i(F), Q, m, F\} = (1, \dots, 1))\}$ .

**D.4.**  $F \in P[m, E] \equiv \cdot^1 (Q_n)\{(1 \leq n \leq 2^{cm^c}) \rightarrow (F \in P(Q_n, m, E))\}$ .

**D.5.**  $F \in P | E | \equiv (\exists m)\{(m \geq n(F)) \wedge (F \in P[m, E])\}$ .

**D.6.**  $E \in P \equiv E \in P | E |$ .

The meaning of the above definitions is analogical to the given in [3] and is explained in [5].

**T.1.** If  $E$  is a thesis, then  $E \in P$ .

The proof of T.1. is inductive on the length of the formal proof of  $E$ , see [3], and is given in [5].

If we replace **D.3.** by:

**D.3'.**  $F \in P(Q, m, E) \equiv J(Q, m, E) \rightarrow V\{i(F), Q, m, F\} = (1, \dots, 1)$ ,

then using Herbrand's proof rules, see [1], we may analogously to [3] prove:

**T.2.** If  $E$  is an alternative of normal forms, then  $E$  is a thesis if and only if  $E \in P$ , see [5].

By an extension of the calculus we mean a first-order functional calculus in which apart of the described signs there are also relations signs  $f_1^1, f_2^1, \dots$  of one argument; in this case the number  $c$  of all relations may be infinite.

Of course, all notations and theorems remain true for the extended calculus; in one we may prove:

**T.3.** A formula  $E$  is a thesis if and only if  $E \in P$ .

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1. Because  $Q$  depends on  $n$ , therefore we write here  $Q = Q_n$ .

We note that in T.5. the number  $c$  which occur in D.4. may be infinite, see [5]<sup>2</sup>; analogical remarks relevant to [3].

T.2-3. prove a new possibility of approximation of the first-order functional calculus by many valued Boolean propositional calculi; in the approximation the quantifier  $\Pi$  is interpreted in T.2. as a finite operator, see (4d).

Some problems connected with T.2-3. we develop in [6]; examples in another paper.

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2. We explain assuming [3]:

To prove the converse theorem to T.1. we prove an analogical theorem to T.2. from [3] in which we assume:

$$R(M) . \equiv . (i) (j) \{ (M/i = M/j) \rightarrow (i = j) \}$$

Then, the theorem holds for all formulas.

But to construct  $M$  with the property  $R(M)$  we use a new sequence of relations  $f_1^1, f_2^1, \dots$ , see [5].