

A THEOREM ON ALTERNATIVES FOR PAIRS OF MATRICES

H. A. ANTOSIEWICZ

The theory of linear inequalities has come into prominence anew in recent years because of its importance in the solution of linear programming problems. In this note we present a simple algebraic proof of an interesting theorem on alternatives for pairs of matrices. This problem was suggested by A. W. Tucker.

Let A and B be matrices, n by m and n by p , respectively, and let x, y, u be column vectors of dimensions m, p, n , respectively.

STATEMENT I. *Either $A'u > 0$, $B'u \geq 0$ for some u or $Ax + By = 0$ for some $x \geq 0$, $y \geq 0$.¹*

STATEMENT II. *Either $A'u \geq 0$, $B'u \geq 0$ for some u or $Ax + By = 0$ for some $x > 0$, $y \geq 0$. [7].*

We shall prove the following theorem.

THEOREM. *Statement I implies, and is implied by, Statement II.*

Note that for the special case when $A = -a$ (column vector) Statement I (or II) reduces to a result of Farkas [2]. If $B = 0$, then Statements I and II are two theorems of Stiemke [6]. More importantly, if the matrix $[B, C, -C]$ is substituted for B , where C is a n by q matrix, and y is replaced by the vector $\begin{pmatrix} y \\ y_1 \\ y_2 \end{pmatrix}$, then Statement I gives the well-known transposition theorem of Motzkin [4, 5]. We refer to [4] for several proofs and further references.

Before proving our theorem, let us make the following preliminary observations. Define the matrix $M = [A, B]$ and the column vector $z = \begin{bmatrix} x \\ y \end{bmatrix}$, and consider the system of equations $Mz = 0$. Assume that the vectors s_1, s_2, \dots, s_k span the linear manifold \mathcal{S} of solutions of this system. Then every solution z can be written in the form $z = S'c$ where $S' = [s_1, s_2, \dots, s_k]$ and c is a k -dimensional (column) vector. Observe that the rows of the matrix M span the orthogonal complement \mathcal{S}^* of \mathcal{S} , that is, every solution of the system $Sz^* = 0$ can be represented as $z^* = M'd$ where d is a n -dimensional (column) vector.

It will be convenient to write $S = [S_1, S_2]$ where S_1 and S_2 are the k by m and k by p matrices, respectively, into which S can be parti-

Received February 3, 1954. This work was performed under a National Bureau of Standards contract with The American University and was sponsored by the Office of Scientific Research, ARDC, USAF.

¹Throughout, transposition is indicated by a dash; also, $x \geq 0$ means $x \geq 0$ with $x = 0$ excluded.

tioned ; accordingly, we introduce two column vectors v, w with m and p components, respectively, and write $z^* = \begin{bmatrix} v \\ w \end{bmatrix}$.

Clearly, the alternatives in each Statement are mutually exclusive as can be seen by multiplying $Ax + By = 0$ on the left by u' . To prove the theorem suppose, at first, that $A'u \geq 0, B'u \geq 0$ for no u and $Ax + By = 0$ has no solution $x > 0, y \geq 0$. Then there exists no c such that

$$S_1'c > 0, \quad S_2'c \geq 0.$$

Hence, by Statement I, the system $S_1v + S_2w = 0$ must be satisfied for some $v \geq 0, w \geq 0$. Since every solution of

$$Sz^* \equiv S_1v + S_2w = 0$$

is of the form $z^* = M'd$, there must exist a vector d such that $A'd \geq 0, B'd \geq 0$, which is a contradiction. Thus Statement I implies Statement II. Conversely, if $A'u > 0, B'u \geq 0$ for no u and $Ax + By = 0$ has no solution $x \geq 0, y \geq 0$, then there exists no c such that $S_1'c \geq 0, S_2'c \geq 0$. Hence, by Statement II, the system $S_1v + S_2w = 0$ must be satisfied for some $v > 0, w \geq 0$, that is, there must exist a vector d such that $A'd > 0, B'd \geq 0$; but this is a contradiction. Thus Statement II implies Statement I.

For applications to linear programming Statements I and II are modified by adjoining in them the inequality $u \geq 0$ to $B'u \geq 0$, that is, by replacing the matrix B by $[B, I]$; in this form they can be used to prove the duality theorem, [1, 3].

REFERENCES

1. G. B. Dantzig, "A proof of the equivalence of the programming problem and the game problem." Cowles Commission Monograph No. 13, "Activity analysis of production and allocation," ed. by T. C. Koopmans, New York 1951.
2. J. Farkas, "Theorie der einfachen Ungleichungen." J. Reine Angew. Math. (Crelle), 124 (1902), 1-27.
3. D. Gale, H. W. Kuhn and A. W. Tucker, "Linear programming and the theory of games." Cowles Commission Monograph No. 13, "Activity analysis of production and allocation," ed. by T. C. Koopmans, New York 1951.
4. T. S. Motzkin, "Beiträge zur Theorie der linearen Ungleichungen." Dissertation, Univ. Basel, 1933.
5. ——— "Two consequences of the transposition theorem on linear inequalities," *Econometrica*, 19 (1951), 184-185.
6. E. Stiemke, "Über positive Lösungen homogener linearer Gleichungen," *Math. Ann.*, 76 (1915), 340-342.
7. A. W. Tucker, "Theorems of alternatives for pairs of matrices." Symposium on Linear Inequalities and programming, USAF and NBS, Washington, D. C., 1951.