CATEGORIES OF V-SETS

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Communicated by Saunders Mac Lane, December 9, 1968

Let V be a partially ordered set. Then a V-set is a function $A: X \to V$ from a set X to V. V is the set of values for A, and X is the carrier of A. If $B: Y \to V$ is another V-set, a morphism $f: A \to B$ is a function $\bar{f}: X \to Y$ such that $A(x) \leq B(\bar{f}(x))$ for each $x \in X$. The category of all V-sets is denoted S(V). The carrier functor $K: S(V) \to S$ assigns X to $A: X \to V$ and $\bar{f}: X \to Y$ to $f: A \to B$, where S is the category of sets. See [2].

If V has one point, S(V) = S. If $V = \{0, 1\}$, where 0 < 1, S(V) is the category of pairs (X, A) of sets, where $A \subseteq X$. If V is the closed unit interval, S(V) is the category of "fuzzy sets", as used by Zadeh and others [1], [5] for problems of pattern recognition and systems theory. When V is a Boolean algebra, V-sets are Boolean-valued sets, as used by Scott and Solovay for independence results in set theory (however, their notion of morphism is different).

If V is complete, S(V) is a pleasant category satisfying all Lawvere's axioms [3] for S except choice, modulo some substitutions of the V-set with carrier 1 and value 0 for the terminal object. In particular,

THEOREM 1. If V is complete, S(V) is complete and cocomplete, has an exponential (i.e., a coadjoint to product) and a "Dedekind-Pierce object" (i.e., an object which looks like the set of integers; see [3]).

Let Poc denote the category of partially ordered classes, and let \mathcal{L} be a subcategory of Poc. Then a category \mathcal{C} is \mathcal{L} -ordered if the power function $\mathcal{O}: |\mathcal{C}| \to Poc$ factors through \mathcal{L} , where $\mathcal{O}(A)$ is the class of all equivalence classes of monics with codomain $A(f \equiv g \text{ if } \exists \text{ an isomorphism } h \text{ such that } fh = g)$. Denote the image of $A \xrightarrow{f} B \text{ by } f(A)$, and the image of the composite $A' \xrightarrow{i} A \xrightarrow{f} B$, where i is monic, by f(A'). Then \mathcal{C} has associative images if it has images such that f(g(A)) = (fg)(A), whenever $A \xrightarrow{g} B \xrightarrow{f} C$. \mathcal{O} can be construed as a functor when \mathcal{C} has associative images. Let $\mathcal{C}L$ denote the category of complete lattices, and call a category C_1 if a coproduct of monics is always monic.

¹ Research supported by Office of Naval Research under contracts Nos. 3656(08) and 222(85), at the University of California at Berkeley.

THEOREM 2. A CL-ordered category with associative images has equalizers, inverse images, unions, intersections, and epic images. If it has coproducts, it is C_1 .

An object P in a category $\mathfrak E$ is monic if every arrow $P \to A$ is monic, and is further atomic if every $P \to A$ is atomic in $\mathfrak P(A)$. P is good if the functor $[P, \]: \mathfrak E \to \mathfrak S$ is noninitial preserving. A union $\bigcup_i A_i$ in $\mathfrak E$ is disjoint if $i \neq j \Rightarrow A_i \cap A_j = \emptyset$, where \emptyset is the initial object. Let CDL be the category of completely distributive lattices, i.e., complete lattices satisfying the law $a \wedge \bigvee_i b_i = \bigvee_i (a \wedge b_i)$. Such lattices V have pseudo-complement operators $*: V \to V$ defined by $a^* = \bigvee_i \{b \mid a \wedge b = 0\}$. Call $V \in |CDL|$ disjointed if for each pair x, y of unequal atoms, $x^* \bigvee_j x = I$, the maximal element of V, and call $\mathfrak E$ disjointedly CDL-ordered if each $\mathfrak P(A) \in |CDL|$ is disjointed.

THEOREM 3. A category C is equivalent to S(V) for some $V \in |CDL|$ if and only if:

- (1) C has an atomic monic good projective generator P;
- (2) C has initial and terminal objects, \(\varphi \) and \(I, \) respectively;
- (3) C has coproducts, which are disjoint unions; and conversely, each disjoint union in C is a coproduct in C;
 - (4) C has associative images;
 - (5) C is disjointedly CDL-ordered; and
 - (6) $P \coprod P$ is not isomorphic to P.

The Axioms (1)-(6) are easily verified for S(V), $V \in |CDL|$. We now sketch the converse, which (surprisingly) makes no use of adjoint functors. Essential use is made of Theorem 2, via Axioms (4) and (5).

Call the elements of [P, A] the *points* of A. We first show the one-pointed objects of $\mathbb C$ are the subobjects of I, except \emptyset ; denote this lattice V. A calculation shows that each $A \in |\mathbb C|$ is a disjoint union $\bigcup_{x \in [P,A]^{z_{\bullet \bullet}}}$, so by Axiom (3), $A = \prod_{x \in [P,A]^{z_{\bullet \bullet}}}$. These facts combine to show that each A is a subobject of $I^{IP,AI}$, the coproduct of I over the index set [P,A]. We then show the arrows $f:A \to B$ in $\mathbb C$ are in 1-1 correspondence with appropriate arrows $\overline{f}:[P,A] \to [P,B]$ in $\mathbb S$. The functor $E:\mathbb C \to \mathbb S(V)$ defined by K(E(A)) = [P,A], $E(A)(x) = x^{**} \in V$, and E(f) = [P,f], is then shown to be full, faithful, and representative.

The addition to Axioms (1)-(6) of either the categorical axiom of choice, or the condition I=P, yields a characterization of S. For finite distributive lattices V, categories of V-sets with *finite* carrier are similarly characterized by all elementary axioms.

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ERRATUM

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Page 125:

Line 3. \tilde{O}^p/\tilde{f}^p should read $\tilde{O}^p/\tilde{f}\tilde{O}^p$.

Line 9. $\alpha(t_m^i) = s_m^i$ should read $\alpha(t_m^i) = s_m^i - \sum_{q=1}^{m-1} a_q s_q^{(i)}$.

Line 10 from bottom. C^N should read Cⁿ.