# PRIMITIVE RECURSIVE COMPUTATIONS 

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1. Definition of a computation.* Using the definition of primitive recursive function found in Kleene, p. 219 [1], we shall define a (primitiverecursive) computation, investigate the mechanics of executing such a computation, and derive upper bounds for the value of the function and for the number of steps required for the computation.

Kleene's definition is: "Each of the following equations and systems of equations (I)-(V) defines a number-theoretic function $\phi$, when $n$ and $m$ are positive integers, $i$ is an integer such that $1 \leqslant i \leqslant n, q$ is a natural number, and $\psi, x_{1}, \ldots, X_{m} \chi$ are given number-theoretic functions of the indicated numbers of variables.

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\begin{equation*}
\phi(x)=x^{\prime} . \tag{I}
\end{equation*}
$$

(II) $\quad \phi\left(x_{1}, \ldots, x_{n}\right)=q$.
(III) $\quad \phi\left(x_{1}, \ldots, x_{n}\right)=x_{i}$.
(IV) $\quad \phi\left(x_{1}, \ldots, x_{n}\right)=\psi\left(\mathrm{X}_{1}\left(x_{1}, \ldots, x_{n}\right), \ldots, \mathrm{X}_{m}\left(x_{1}, \ldots, x_{n}\right)\right)$.
(Va) $\quad\left\{\begin{aligned} \phi(0) & =q, \\ \phi\left(y^{\prime}\right) & =\chi(y, \phi(y)) .\end{aligned}\right.$
(Vb) $\left\{\begin{aligned} \phi\left(0, x_{2}, \ldots, x_{n}\right) & =\psi\left(x_{2}, \ldots, x_{n}\right), \\ \phi\left(y^{\prime}, x_{2}, \ldots, x_{n}\right) & =\mathrm{X}\left(y, \phi\left(y, x_{2}, \ldots, x_{n}\right), x_{2}, \ldots, x_{n}\right) .\end{aligned}\right.$
(( Va ) constitutes the case of ( V ) for $n=1$, and ( Vb ) for $n>1$.) A function is primitive recursive if it is definable by a series of applications of these five operations of definition."

Modifying this definition to permit zero arguments in (II) so that (Va) and ( Vb ) can be combined, we proceed in the obvious way to give a recursive definition of function word, giving in the process a definition of the rank of a function word:
(1) $S$ is a function word of rank 1.
(2) $C_{m}^{n}$ is a function word of rank $n(n \geqslant 0)$.
(3) $U_{m}^{n}$ is a function word of rank $n(n \geqslant 1 ; 1 \leqslant m \leqslant n)$.
(4) If $A^{m}$ is a function word of rank $m$ and if $B_{1}^{n}$, .., $B_{m}^{n}$ are function words of rank $n$, then $\mathbf{S}_{m}^{n} A^{m} B_{1}^{n} \ldots B_{m}^{n}$ is a function word of rank $n$ ( $m \geqslant 1, n \geqslant 1$ ).

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