

A STRONG COMPLETENESS THEOREM FOR 3-VALUED LOGIC: PART II

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Proof¹ was given in [1] that $\mathbf{SC}_3$, the 3-valued sentential calculus, has a strongly complete axiomatization. Pushing our investigation one step further,² we obtain here a like result about $\mathbf{QC}_3$, the 3-valued quantificational calculus of order one.³

1 The primitive signs of $\mathbf{QC}_3$ are

- (a) ' $\sim$ ', ' $\supset$ ', ' $\forall$ ', ' $(, ')$ ', and ' $,$ ',
- (b) a denumerable infinity of individual variables, to be referred to by means of ' X ',⁴
- (c) a denumerable infinity of individual parameters, to be referred to by means of ' X ',⁵ and
- (d) for each d from 0 on, a denumerable infinity of predicate parameters of degree d , to be referred to by means of ' F^d '.⁶

We presume the variables in (b), the parameters in (c), and the parameters in (d) to be alphabetically ordered; and we take the alphabetically first parameter of degree d in (d) to be ' p '.

The atomic wffs of $\mathbf{QC}_3$ are all formulas of the sort $F^d(X_1, X_2, \dots, X_d)$, where F^d is a predicate parameter of degree d ($d \geq 0$) and $X_1, X_2, \dots$, and X_d are individual parameters. The wffs of $\mathbf{QC}_3$ (presumed at one point below to be alphabetically ordered) are the atomic wffs just defined, plus all formulas of the sorts (i) $\sim A$, where A is well-formed, (ii) $(A \supset B)$, where A and B are well-formed, and (iii) $(\forall X)A$, where—for some individual parameter X —the result $A(X/X)$ of replacing X everywhere in A by X is well-formed.⁷ The length $\mathcal{L}(A)$ of an atomic wff is 1; the length $\mathcal{L}(\sim A)$ of a negation $\sim A$ is $\mathcal{L}(A) + 1$; the length $\mathcal{L}((A \supset B))$ of a conditional $(A \supset B)$ is $\mathcal{L}(A) + \mathcal{L}(B) + 1$; and the length $\mathcal{L}((\forall X)A)$ of a quantification $(\forall X)A$ is $\mathcal{L}(A(X/X)) + 1$, where X is the alphabetically earliest individual parameter of $\mathbf{QC}_3$. We avail ourselves of the following ten abbreviations:

$$\begin{aligned}
 \text{'f'} &=_{df} \text{'}\sim(p \supset p)\text{' } \\
 (A \vee B) &=_{df} ((A \supset B) \supset B)^8 \\
 (A \& B) &=_{df} \sim(\sim A \vee \sim B) \\
 (A \equiv B) &=_{df} ((A \supset B) \& (B \supset A))
 \end{aligned}$$