A MORE RELEVANT RELEVANCE LOGIC

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Relevant implication or entailment is designed to convey the notion of "logical consequence". While $A \supset B$, where \supset is the weak implication usually specified by a truth table, tells us that either A is false or B is is true, $A \to B$, where \to is entailment, tells us that A is actually used, and perhaps necessary, in the proof of B. In this paper we show that, in a certain sense, the entailment systems of Anderson and Belnap ([1]) are still not fully relevant and we describe a new system which is at least more so.

The simplest notion of relevance can be expressed in terms of the following deduction theorem:

If there is a proof of B using all of
$$A_1, \ldots, A_n$$
, then $\vdash A_1 \rightarrow A_2 \rightarrow \ldots \rightarrow A_n \rightarrow B$.

(The A_i s, B, and all capital Roman letters used below range over (well formed) formulas as in [1]).

This deduction theorem together with modus ponens or $\rightarrow E$ (\rightarrow elimination) is equivalent to the system R_{\rightarrow} of [1]. This system can be axiomatized as follows:

$$\begin{array}{lll} \mathbf{R}_{\rightarrow}\mathbf{1} & \vdash A \rightarrow A \\ \mathbf{R}_{\rightarrow}\mathbf{2} & \vdash A \rightarrow B \rightarrow . \ B \rightarrow C \rightarrow . \ A \rightarrow C \\ \mathbf{R}_{\rightarrow}\mathbf{3} & \vdash (A \rightarrow . \ B \rightarrow C) \rightarrow . \ B \rightarrow . \ A \rightarrow C \\ \mathbf{R}_{\rightarrow}\mathbf{4'} & \vdash (A \rightarrow . \ B \rightarrow C) \rightarrow . \ A \rightarrow B \rightarrow . \ A \rightarrow C^{1} \end{array}$$

The instance

$$\vdash A \to A \to A \to A \tag{1}$$

of R_1 and R_2 lead directly to

$$\vdash A \rightarrow . (A \rightarrow A) \rightarrow A.$$
 (2)

within which the $A \rightarrow A$ still seems to be irrelevant. Anderson and Belnap also claim (2) to be irrelevant, but perhaps for other reasons. They