

# BETH'S TABLEAUX FOR RELEVANT LOGIC

J. F. PABION

**1 Sequents—Countermodels** For the concept of *relevant-model-structure* (r.m.s.) we\* refer to [2]. Giving a r.m.s.  $\mathfrak{M} = \langle K, 0, R, * \rangle$  satisfying postulates  $P_1$ – $P_6$  of [2], a “forcing relation” is defined between  $K$  and the set of atomic sentences defined in a given domain  $D$  such that:

$$\alpha \Vdash A \text{ and } R(0, \alpha, \beta) \Rightarrow \beta \Vdash A.$$

Then the forcing relation is extended according to usual stipulations:

$$\begin{aligned} \alpha \Vdash A \vee B &\Leftrightarrow \alpha \Vdash A \text{ or } \alpha \Vdash B \\ \alpha \Vdash A \wedge B &\Leftrightarrow \alpha \Vdash A \text{ and } \alpha \Vdash B \\ \alpha \Vdash A \rightarrow B &\Leftrightarrow \forall \beta, \gamma, \beta \Vdash A \text{ and } R(\alpha, \beta, \gamma) \Rightarrow \gamma \Vdash B \\ \alpha \Vdash \bar{A} &\Leftrightarrow \alpha^* \nVdash A \\ \alpha \Vdash \forall x F[x] &\Leftrightarrow \forall a \in D, \alpha \Vdash F[a]. \end{aligned}$$

A *signed formula* is an expression of one of the forms:

$$+A \quad -A$$

where  $A$  is a formula. A *sequent* is a finite sequence of signed sentences.

A *countermodel* of a sequent is given by:

$$\text{A r.m.s. } \mathfrak{M} = \langle K, 0, R, * \rangle.$$

A member  $\alpha$  of  $K$  such that:

$$\begin{aligned} \text{If } +A \text{ is in the sequent, then } \alpha \Vdash A \\ \text{If } -A \text{ is in the sequent, then } \alpha \nVdash A. \end{aligned}$$

$\langle \mathfrak{M}, \alpha \rangle$  is a *normal countermodel* if  $\alpha = 0$  and  $0^* = 0$ . The sequent is said to be *valid* if it has no countermodel and *normally valid* if it has no normal countermodel.

---

\*The author is grateful to professors Belnap, Dunn, and Meyer for having focalized his attention on typical mistakes contained in some tentative proofs.